

By Lemma 5 the statement forced is  $\Pi_2^1$  in constructible parameters and so must be satisfied by every non-constructible path through  $P$ . Since there is such a path, this contradicts our original assumption that  $<$  is a linear ordering, and establishes the claim.

It is easy to see that the only elements of  $B$  invariant under all automorphisms of  $B$  are 0 and 1. From this it follows that in  $L(S)$  all definable sets are constructible [11], [6, Theorem 6.8]. However it must also be valid that the first non-constructible element in the ordering  $<$  is definable and non-constructible.

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## On longest paths in connected graphs\*

by

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**Abstract.** It is shown that a connected graph of order  $p \geq 4$  contains a path of length  $k$ , where  $1 \leq k \leq p-1$ , if for every integer  $j$  with  $1 \leq j < k/2$ , the number of vertices of degree not exceeding  $j$  is less than  $j$ . Furthermore, a tree of order  $p \geq 4$  has diameter at least  $k$ , where  $3 \leq k \leq p-1$ , if the number of vertices of degree one is less than  $\{2(p-1)/(k-1)\}$ .

A *hamiltonian cycle (path)* in a graph  $G$  is a cycle (path) containing every vertex of  $G$ . Pósa [1] proved that if  $G$  is a graph of order  $p \geq 3$  such that for every integer  $j$  with  $1 \leq j < p/2$ , the number of vertices of degree not exceeding  $j$  is less than  $j$ , then  $G$  contains a hamiltonian cycle. In this article, we establish an analogous result for graphs with hamiltonian paths and in fact for graphs containing paths of any specified length.

**THEOREM 1.** *Let  $G$  be a connected graph of order  $p \geq 4$ . Then  $G$  contains a path of length  $k$  ( $1 \leq k \leq p-1$ ) if for every integer  $j$  with  $1 \leq j < k/2$  the number of vertices of degree not exceeding  $j$  is less than  $j$ .*

**Proof.** Since  $G$  is connected and  $p \geq 4$ , the theorem is true for  $k = 1$  and  $k = 2$ . Henceforth we assume  $k \geq 3$ . Suppose the length of a longest path in  $G$  is  $n$  where  $2 \leq n < k$ . If  $P$  is a longest path in  $G$ , let  $S_P$  denote  $\deg u + \deg v$ , where  $u$  and  $v$  are the endvertices of  $P$ . Among all longest paths in  $G$ , choose  $P$ :  $u_0, u_1, \dots, u_n$  so that  $S_P$  is maximum. Suppose  $\deg u_0 \leq \deg u_n$ .

Since  $P$  is a longest path, each of  $u_0$  and  $u_n$  is adjacent only to vertices of  $P$ . If  $u_i u_n \in E(G)$ ,  $0 \leq i \leq n-1$ , then  $u_0 u_{i+1} \notin E(G)$ ; for otherwise the cycle

$$C: u_0, u_1, \dots, u_i, u_n, u_{n-1}, \dots, u_{i+1}, u_0$$

of length  $n+1$  is present in  $G$ . The cycle  $C$  cannot contain all vertices of  $G$  since  $n+1 < p$ . Since  $G$  is connected, there exists a vertex  $w$  not on  $C$  adjacent to a vertex of  $C$ ; however this implies  $G$  contains a path of length  $n+1$  which is impossible. Hence for each vertex of  $\{u_0, u_1, \dots, u_{n-1}\}$

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adjacent to  $u_n$ , there is a vertex of  $\{u_1, u_2, \dots, u_n\}$  not adjacent to  $u_0$ . Thus  $\deg u_0 \leq n - \deg u_n$ , so that  $\deg u_0 + \deg u_n \leq n$ . Since  $\deg u_0 \leq \deg u_n$  and  $n < k$  we obtain  $\deg u_0 < k/2$ , so, by hypothesis, there are fewer than  $\deg u_0$  vertices having degree not exceeding  $\deg u_0$ . However, if  $u_0 u_i \in E(G)$ ,  $1 \leq i \leq n$ , then

$$u_{i-1}, u_{i-2}, \dots, u_0, u_i, u_{i+1}, \dots, u_n$$

is a longest path in  $G$  and from the manner in which  $u_0$  and  $u_n$  were chosen,  $\deg u_{i-1} \leq \deg u_0$ . Thus there are at least  $\deg u_0$  vertices having degree not exceeding  $\deg u_0$ . This presents a contradiction so that  $k \leq n$ . Thus  $G$  contains a path of length  $k$ .

The case  $k = p - 1$  gives a sufficient condition for a graph to possess a hamiltonian path. However, we can obtain a slightly stronger result by employing the previously stated theorem of Pósa. Suppose  $G$  is a graph of order  $p \geq 3$  such that for every integer  $j$  with  $0 \leq j < (p - 1)/2$ , the number of vertices of degree not exceeding  $j$  is less than  $j + 1$ . We construct a graph  $G^*$  by adding a new vertex  $v$  and new edges joining  $v$  to all vertices of  $G$ . The graph  $G$  has a hamiltonian path if and only if  $G^*$  has a hamiltonian cycle. However,  $G^*$  satisfies the hypothesis of Pósa's theorem (with  $p + 1$  in place of  $p$ ) so that  $G^*$  is hamiltonian.

The results of Theorem 1 are, in general, not best possible. However, the result stated in the preceding paragraph is best possible. For example, if  $p \geq 4$ , let  $G$  be a connected graph with one cutvertex and the three blocks  $K_2, K_{n+1}$ , and  $K_{p-n-1}$ , where  $1 \leq n < (p - 1)/2$ . Then  $G$  satisfies all the hypotheses with one exception:  $G$  contains at least  $n + 1$  vertices of degree not exceeding  $n$ . Furthermore, the graph  $G$  fails to contain a hamiltonian path.

The simplest connected graphs are the trees. In this case, one can give a sufficient condition for lengths of longest paths in terms of the number of vertices of degree one (endvertices) only. Since there is a unique path between every pair of vertices in a tree, the result can be expressed in terms of its diameter (the length of a longest path for a tree). We denote the diameter of a graph  $G$  by  $\text{diam} G$ .

We begin with two lemmas.

**LEMMA 1.** *If a nontrivial tree  $T$  has  $m$  endvertices, where  $m$  is even, then there exist paths  $P_1, P_2, \dots, P_{m/2}$  satisfying:*

- (1) *Each endvertex of  $T$  is an endvertex of exactly one of the paths  $P_i$ ,  $1 \leq i \leq m/2$ ,*
- (2) 
$$\bigcup_{j=1}^{m/2} V(P_j) = V(T),$$
- (3) 
$$V(P_i) \cap \left( \bigcup_{j=1}^{i-1} V(P_j) \right) \neq \emptyset \quad \text{for all } i \text{ with } 2 \leq i \leq m/2.$$

**Proof.** We employ induction on  $m$ . For  $m = 2$ ,  $T$  is a path and the result desired is satisfied by  $T$  itself. Assume the result holds for all trees with  $m \geq 2$  endvertices, and let  $T$  be a tree with  $m + 2$  endvertices. With each endvertex of  $T$ , we associate a path of  $T$  in the following manner. Let  $u$  be an arbitrary endvertex of  $T$ . Now,  $\deg u = 1$ , where  $\dot{u}$  is adjacent to, say,  $u_1$ . If  $\deg u_1 \geq 3$ , let  $P_u$  be the trivial path:  $u$ . Otherwise,  $\deg u_1 = 2$  and  $u_1$  is adjacent to  $u$  and, say,  $u_2$ . If  $\deg u_2 \geq 3$ , let  $P_u$  be the path:  $u, u_1$ . Otherwise,  $\deg u_2 = 2$  and  $u_2$  is adjacent to  $u_1$  and, say,  $u_3$ . Continue in this fashion to obtain the path  $P_u$ . Since  $m + 2 \geq 4$ , we can choose endvertices  $u$  and  $v$  of  $T$  so that  $T^* = T - V(P_u) - V(P_v)$  is a tree with  $m$  endvertices, each of which is an endvertex of  $T$ . By the induction hypothesis,  $T^*$  contains a collection of paths  $P_1, P_2, \dots, P_{m/2}$  satisfying conditions (1)-(3). Let  $u'$  and  $v'$  be the final vertices of  $P_u$  and  $P_v$ , respectively. Then there exists a  $z-w$  path  $P$  in  $T^*$  where  $u'z, wv' \in E(T)$ . If we define  $P_{(m+2)/2}$  to be the path  $P_u$  followed successively by  $u'z, P, wv'$ , and  $P_v$ , the result follows.

**LEMMA 2.** *If a nontrivial tree  $T$  has  $m$  endvertices, where  $m$  is odd, and  $\text{diam} T = k$ , then there exist paths  $P_1, P_2, \dots, P_{(m-1)/2}, P_{(m+1)/2}$  satisfying:*

- (1)  $P_{(m+1)/2}$  is of the form  $P_u$  (using the notation of the previous lemma) for some endvertex  $u$  of  $T$  and the length of  $P_u$  does not exceed  $(k/2) - 1$ ,
- (2) Each endvertex of  $T$  other than  $u$  is an endvertex of exactly one of the  $P_i$ 's for  $1 \leq i \leq (m - 1)/2$ ,
- (3) 
$$V(P_i) \cap \left( \bigcup_{j=1}^{i-1} V(P_j) \right) \neq \emptyset \quad \text{for all } i \text{ with } 2 \leq i \leq (m - 1)/2.$$

**Proof.** Since  $\text{diam} T = k$ , the path  $P_{(m+1)/2}$  clearly exists. Let  $T^* = T - V(P_{(m+1)/2})$ . Then  $T^*$  is a tree with  $m - 1$  endvertices, each of which is an endvertex of  $T$ . By applying the previous lemma to  $T^*$ , we obtain the desired result.

**THEOREM 2.** *Let  $T$  be a tree of order  $p \geq 4$  and let  $k$  be a fixed integer with  $3 \leq k \leq p - 1$ . If  $\text{diam} T < k$ , then  $T$  has at least  $\{2(p - 1)/(k - 1)\}$  endvertices.*

**Proof.** Let  $m$  denote the number of endvertices of  $T$ . We consider two cases depending on the parity of  $m$ .

**Case 1.**  $m$  is even. By Lemma 1, there exist paths  $P_1, P_2, \dots, P_{m/2}$  with the three required properties. Since  $\text{diam} T \leq k - 1$ , we have  $|V(P_i)| \leq k$  for  $i = 1, 2, \dots, m/2$ . We wish to find an upper bound for  $p$ . Since  $|V(P_1)| \leq k$ ,  $|V(P_2)| \leq k$ , and  $V(P_2) \cap V(P_1) \neq \emptyset$ , we have

$$|V(P_1) \cup V(P_2)| \leq k + (k - 1).$$

Moreover, since  $V(P_3) \cap (\bigcup_{j=1}^2 V(P_j)) \neq \emptyset$  and  $|V(P_3)| \leq k$ , it follows that  $|V(P_1) \cup V(P_2) \cup V(P_3)| \leq k + (k-1) + (k-1)$ . Continuing in this fashion, we obtain  $|\bigcup_{j=1}^{m/2} V(P_j)| \leq (m/2)(k-1) + 1$ . Since  $V(T) = \bigcup_{j=1}^{m/2} V(P_j)$ , we have  $p \leq (m/2)(k-1) + 1$  so that  $m \geq 2(p-1)/(k-1)$ . Thus  $m \geq \{2(p-1)/(k-1)\}$  since  $m$  is integral.

Case 2.  $m$  is odd. By Lemma 2, there exist paths  $P_1, P_2, \dots, P_{(m+1)/2}$  with the four required properties. Since  $\text{diam } T \leq k-1$ ,  $|V(P_i)| \leq k$  for  $i = 1, 2, \dots, (m-1)/2$  and  $|V(P_{(m+1)/2})| \leq (k-1)/2$ . As above, we obtain

$$|\bigcup_{j=1}^{(m-1)/2} V(P_j)| \leq \left(\frac{m-1}{2}\right)(k-1) + 1$$

and hence

$$|\bigcup_{j=1}^{(m+1)/2} V(P_j)| \leq \left(\frac{m-1}{2}\right)(k-1) + 1 + \frac{(k-1)}{2}.$$

Since  $V(T) = \bigcup_{j=1}^{(m+1)/2} V(P_j)$ , we have  $p \leq (m/2)(k-1) + 1$  so that

$$m \geq 2(p-1)/(k-1).$$

Thus  $m \geq \{2(p-1)/(k-1)\}$  since  $m$  is integral.

The previous theorem may be restated as follows.

**THEOREM 2'.** *Let  $T$  be a tree of order  $p \geq 4$  and let  $k$  be a fixed integer with  $3 \leq k \leq p-1$ . If  $T$  has fewer than  $\{2(p-1)/(k-1)\}$  endvertices, then  $\text{diam } T \geq k$ .*

This result is best possible in the following sense. Given an odd integer  $k \geq 3$  and integer  $m \geq 2$ , there exists an integer  $p \geq 4$  and tree  $T$  of order  $p$  and diameter  $k-1$  with  $3 \leq k \leq p-1$ , such that  $T$  has  $m$  endvertices, where  $m = \{2(p-1)/(k-1)\}$ . The tree obtained by replacing each edge of the (star) graph  $K(1, m)$  with a path of length  $(k-1)/2$  serves as such an example.

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