

Representation of involuted semigroups by binary relations

by

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Abstract. We give necessary and sufficient conditions for an (ordered) involuted semigroup to be isomorphic with an (ordered) involuted semigroup of binary relations in which the involution is the conversion of relations (and the order is the set-theoretical inclusion). The same problem is solved for representations of (ordered) semigrouds (or semiheaps) by binary relations. The conditions of representability are universal Horn formulas.

An involuted semigroup is an algebra $A=(A;\cdot,^{-1})$ with a binary multiplication \cdot and unary involution $^{-1}$ subject to the following three identities:

$$(1) (xy)z = x(yz),$$

$$(2) (x^{-1})^{-1} = x ,$$

$$(3) (xy)^{-1} = y^{-1}x^{-1}.$$

A subalgebra of an involuted semigroup A is called an *involuted sub-semigroup* of A.

Historically, the first example of involuted semigroups was that of groups. Every group is isomorphic to an involuted semigroup of permutations of a set (the Cayley theorem). Here a permutation of a set X is a bijective self-mapping of X, i.e. a one-to-one and full binary relation on X. The multiplication and involution of permutations have their usual meaning of composition and conversion.

If in the definition of permutations we do not demand that they should be full, i.e. everywhere defined, we obtain involuted semigroups of partial one-to-one transformations of a set. Such involuted semigroups are inverse semigroups and every inverse semigroup is isomorphic to an involuted semigroup of one-to-one partial transformations of a set (the Vagner-Preston theorem).

If in the definition of permutations we do not demand that they should be one-to-one, we obtain full multi-valued transformations, i.e. full



binary relations of a set (in effect, since involuted semigroups are closed under involution, an involuted semigroup of binary relations contains converses of all the relations belonging to it; if at least one of the relations is not one-to-one, the involuted semigroup contains a multi-valued transformation). Those involuted semigroups which are isomorphic to involuted semigroups of full binary relations form an extensive class of algebras containing, in particular, all groups and all inverse semigroups. An axiomatization of this class was found in [12]. In particular, an involuted semigroup is an inverse semigroup if and only if it is isomorphic to an involuted semigroup of full binary relations and satisfies the identity [12] and [13]

$$(4) xx^{-1}x = x.$$

Now, if in the definition of permutations we do not demand either one-to-oneness or fullness, we obtain arbitrary binary relations. An *involuted* semigroup of binary relations is an involuted semigroup of the form $(F; \circ, ^{-1})$ where F is a nonempty set of binary relations on a set, \circ is the relative multiplication and $^{-1}$ is the usual conversion of binary relations.

Thus involuted semigroups of binary relations are, in a sense, an ultimate generalization of permutation groups. The main problem solved in this paper is an explicit abstract characterization of involuted semigroups of binary relations. In other words, we give a system of axioms necessary and sufficient in order for an abstract involuted semigroup to be isomorphic to an involuted semigroup of binary relations. This problem is similar to the problem of characterizing relation algebras in the sense of A. Tarski [17] or relation algebras in the sense of B. Jónsson [7]. An involuted semigroup isomorphic to an involuted semigroup of binary relations is called representable. It is known that the class of representable involuted semigroups forms a quasivariety (i.e. it can be characterized by a system of elementary axioms each of which is a conditional identity) and not every involuted semigroup is representable [14]. It is shown in the present paper that the class of all representable involuted semigroups does not form a variety of algebras, i.e. cannot be characterized by a set of identities.

As a first step we solve the problem of abstract characterization of representable ordered involuted semigroups. An ordered involuted semigroup is an algebraic system of the form $(A; \cdot, \cdot^{-1}, \leq)$ where $(A; \cdot, \cdot^{-1})$ is an involuted semigroup and $(A; \leq)$ is a (partially) ordered set, where the order relation \leq is stable under both operations of the involuted semigroup, i.e. satisfies the axioms:

$$(5) x_1 \leqslant x_2 \& y_1 \leqslant y_2 \Rightarrow x_1 y_1 \leqslant x_2 y_2,$$

$$(6) x \leqslant y \Rightarrow x^{-1} \leqslant y^{-1}.$$

Since every set of binary relations is ordered by the set-theoretical inclusion of relations, we may consider ordered involuted semigroups of binary relations $(F; \circ, \neg^{-1}, \subset)$; clearly, they satisfy conditions (5) and (6). An ordered involuted semigroup is called representable if it is isomorphic to an ordered involuted semigroup of binary relations.

The method used to solve the problem of characterization of representable involuted semigroups permits us to solve a similar problem of characterizing representable semigrouds. This problem is considered by many authors as the central problem of the theory of semigrouds.

A semigroud (in other terminology, a semiheap) is an algebra (A; []) with a ternary operation $(a_1, a_2, a_3) \mapsto [a_1 a_2 a_3]$ satisfying the following identities of skew associativity

(7)
$$|[x_1x_2x_3]x_4x_5| = |[x_1[x_4x_3x_2]x_5] = |[x_1x_2[x_3x_4x_5]].$$

Semigrouds are closely connected with involuted semigroups. In fact, if $(A; \cdot, \cdot^{-1})$ is an involuted semigroup and $[a_1a_2a_3] = a_1a_2^{-1}a_3$, then $(A; [\cdot])$ is a semigroud and every semigroud can be isomorphically embedded into a semigroud associated with an involuted semigroup [2]. An ordered semigroud is an algebraic system of the form $(A; [\cdot], \leq)$ where $(A; [\cdot])$ is a semigroud, $(A; \leq)$ is a (partially) ordered set and the order relation \leq is stable under the ternary multiplication $[\cdot]$:

(8)
$$x_1 \leqslant x_2 \& y_1 \leqslant y_2 \& z_1 \leqslant z_2 \Rightarrow [x_1 y_1 z_1] \leqslant [x_2 y_2 z_2].$$

Let ϱ , σ , $\tau \subset X \times Y$ be binary relations between the elements of two sets X and Y. Define $[\varrho\sigma\tau] = \tau \circ \sigma^{-1} \circ \varrho$ (the factors in the product of binary relations are written from right to left). Then $[\varrho\sigma\tau] \subset X \times Y$ and the set of all binary relations between the elements of X and Y is a semigroud under the ternary multiplication $[\]$. Subsemigrouds of such semigrouds (for arbitrary X and Y) are called semigrouds of binary relations. If $(F; [\])$ is a semigroud of binary relations, then $(F; [\], \subset)$ is an ordered semigroud of binary relations (here \subset is the set-theoretical inclusion relation on F). An (ordered) semigroud is called representable if it is isomorphic to an (ordered) semigroud of binary relations. It is known [14] that the class of all representable semigrouds forms a quasivariety and not every semigroud is representable. The problems of characterizing representable involuted semigroups and semigrouds are long-standing ones (see [2], where the concept of a semigroud was introduced for the first time, and Problem T21 in [19]).

The problems of abstract characterizations of representable (ordered) involuted semigroups and semigrouds are typical problems on relation algebras in the sense of [10] and [11]. Representations of semigrouds by binary relations have been considered in [4], where a system of axioms

for representable semigrouds has been found. However, these axioms are sentences in the third order predicate calculus. By the fundamental theorem on relation algebras, [10] and [11], the class of representable semigrouds is elementarily axiomatizable. We find a system of elementary axioms for this class.

A system of elementary axioms for representable involuted semigroups has been found by R. McKenzie [8] (never published). However, these axioms have a very complicated form and are not conditional identities. Our axioms have a much simpler form. Elementary axioms for representable ordered involuted semigroups have been found by D. A. Bredihin [1]. These axioms are similar to ours; however, our axioms are visualized much more easily.

Abstract characterizations for (ordered) involuted semigroups of special binary relations have been found in [5], [12], [13] and [16].

Abstract characterizations of (ordered) involuted semigroups of all binary relations on a set and of semigrouds of all binary relations between the elements of two sets have been found in [3], [6], [9] and [14] (in the case of semigroups) and in [14] (in the case of semigrouds). A related problem has been solved in [15].

The following concept of net is of central importance in this paper. A net of order n is a finite oriented graph M_n having n+1 vertices (which we identify with the natural numbers $\{0,1,\ldots,n\}$), having no loops, having an arrow (0,1) (here (i,j) denotes an arrow, i.e. an oriented edge, leading from i to j) but not an arrow (1,0), and such that for every vertex i>1 there exists a uniquely determined arrow (p_i,i) leading from $\{0,1,\ldots,i-1\}$ to i and a uniquely determined arrow (i,q_i) leading from i to $\{0,1,\ldots,i-1\}$ (thus, $p_i < i$ and $q_i < i$).

Since nets are extensively used in what follows, we will consider this concept in more detail.

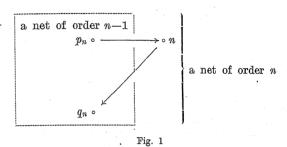
First of all, it is clear that if M_n is a net and $0 < k \le n$, then the set $\{0, 1, ..., k\}$ of vertices and the set of all arrows of M_n between these vertices form a net M_k of order k. A net M is called a *subnet* of a net M_n if $M = M_k$ for some k, $0 < k \le n$. Thus M_n has precisely n subnets, namely, $M_1, M_2, ..., M_n$.

Any net M_n of order n has exactly 2n-1 arrows. In effect, the subnet M_1 of M_n has one arrow, by definition. Suppose the subnet M_{i-1} has 2i-3 arrows. The subnet M_i is formed from M_{i-1} by adding a new vertex i and two new arrows (p_i, i) and (i, q_i) . Thus M_i contains 2i-3+2=2i-1 arrows. By induction, M_n contains 2n-1 arrows.

We shall number consecutively the arrows of any net of order n. The arrow (0, 1) has number 0. The arrows (p_i, i) and (i, q_i) have numbers 2i-3 and 2i-2 respectively. Thus every arrow is numbered by a single number from 0 to 2n-2. We identify the arrows with their numbers;

to distinguish them from vertices, arrows are written as bold face numbers. E.g. $(p_i, i) = 2i - 3$.

All nets can be constructed by a very simple procedure. The only existing net M_1 of order 1 is just two vertices 0 and 1 and an arrow 0 = (0, 1). Any net of order n is obtained from a net M_{n-1} of order n-1 by adding a new vertex n and two new arrows $2n-3 = (p_n, n)$ and $2n-2 = (n, q_n)$ where p_n and q_n are some (possibly coinciding) fixed vertices of M_{n-1} . This may be visualized in Fig. 1.



Thus, a net of order n can be constructed from a net M_{n-1} of order n-1 in n^2 different ways. It follows that there exist $1 \cdot 2^2 \cdot \dots \cdot n^2 = (n!)^2$ different nets of order n.

If (i,j) is an arrow, then $(i,j)^{-1}$ is called the reversed arrow. The reversed net M_n^{-1} is obtained from M_n if all the arrows of M_n are reversed. $(i,j)^{-1}$ leads from j to i.

An edge of M_n is either an arrow of M_n or an arrow of M_n^{-1} (i.e. an edge is either an arrow or a reversed arrow of M_n). A chain in M_n is a nonempty finite sequence of edges such that the end of each edge coincides with the beginning of the next edge.

Let $A = (A; \cdot, \neg^{-1}, \leqslant)$ be an ordered involuted semigroup and let M_n be a net of order n. A metric on M_n (more precisely, an A-metric on M_n) is any mapping $d: \{0, 1, ..., 2n-2\} \rightarrow A$ of the set of arrows of M_n into A.

Define $d(i^{-1}) = d(i)^{-1}$ for every arrow i of M_n . If $\mu = (e_1, e_2, ..., e_m)$ is a chain in M_n consisting of edges e_i , write $d(\mu) = d(e_1)d(e_2)...d(e_m)$, where the right-hand side product is taken inside A. Then $d(\mu)$ is called the length of μ .

An A-metric on M_n is called *proper* if for any k, 0 < k < n, there exists a chain μ_k in M_k leading from p_{k+1} to q_{k+1} whose length is shorter than that of the chain (2k-1, 2k), i.e. $d(\mu_k) \leq d(2k-1)d(2k)$.

Proper metrics on M_n can be constructed inductively in a simple way. We can define d consecutively on subnets $M_1, M_2, ..., M_n$. A proper metric on M_1 is merely a mapping of $\{0\}$ into A, i.e. to define it we should



pick out an element $a_0 \in A$ and suppose $d(0) = a_0$. Suppose d is defined on M_k . Choose a chain μ_k in M_k leading from p_{k+1} to q_{k+1} and pick out two elements $a_{2k-1}, a_{2k} \in A$ such that $d(\mu_k) \leq a_{2k-1} a_{2k}$. Then define $d(2k-1) = a_{2k-1}, d(2k) = a_{2k}$. The resulting d is a proper metric on M_{k+1} . Thus we can construct a proper metric on M_n and every proper metric on M_n may be constructed in this fashion. We have supposed that for every $a \in A$ there exist $b, c \in A$ such that $a \leq bc$. However, as we shall see later, this is always the case in those ordered involuted semigroups in which we are interested.

A proper A-metric d on M_n is called regular if for every chain μ in M_n leading from 0 to 1 $d(0) \leq d(\mu)$.

Theorem 1. An ordered involuted semigroup A is representable if and only if every proper A-metric on every net is regular.

Proof. Necessity. Suppose d is a proper F-metric on a net M_n of order n for an ordered involuted semigroup F of binary relations on a set X. Let $(x_0, x_1) \in d(\mathbf{0})$ for $x_0, x_1 \in X$.

We are going to define a sequence $x_0, x_1, ..., x_n$ of elements of X such that if (i, j) is an arrow of M_n , then $(x_i, x_j) \in d((i, j))$.

The first two terms of such a sequence have already been given. Suppose $x_0, ..., x_k$ have already been defined $(0 < k \le n)$. By supposition, there exists a chain μ_k leading from p_{k+1} to q_{k+1} and such that $d(\mu_k) \subset d(2k-1)d(2k)$. By the supposition of induction, $(x_i, x_j) \in d((i, j))$ for $0 < i, j \le k$. Therefore, $(x_j, x_i) \in d((i, j))^{-1} = d((i, j)^{-1})$. By the definitions of a chain and the relative multiplication, $(x_i, x_j) \in d(v)$, where v is any chain in M_k leading from i to j. In particular, $(x_{p_{k+1}}, x_{q_{k+1}}) \in d(\mu_k) \subset d(2k-1)d(2k)$. (Notice that here we do not use \circ to denote relative multiplication and the factors in the relative product of two binary relations are written from left to right and not from right to left as we do when \circ is used). It follows that $(x_{p_{k+1}}, x_{k+1}) \in d(2k-1)$ and $(x_{k+1}, x_{q_{k+1}}) \in d(2k)$ for some $x_{k+1} \in X$. Choose such x_{k+1} . Clearly, now $(x_i, x_j) \in d(i, j)$ for $0 < i, j \le k+1$. By induction, the whole sequence $x_0, ..., x_n$ can be constructed.

Now suppose μ is a chain in M_n leading from 0 to 1. As we have just seen, this implies $(x_0, x_1) \in d(\mu)$. Thus $(x_0, x_1) \in d(0)$ implies $(x_0, x_1) \in d(\mu)$ for all $x_0, x_1 \in X$, i.e. $d(0) \subseteq d(\mu)$. Therefore, the metric d is regular.

Now if A is a representable ordered involuted semigroup, then every proper A-metric is regular since this is true for all proper metrics with values in the ordered involuted semigroup of binary relations which is an isomorphic image of A.

Sufficiency. Let K denote the class of all representable ordered involuted semigroups. Suppose A is an ordered involuted semigroup and every proper A-metric on every net is regular. Then this is true for all

proper B-metrics where B is any ordered involuted subsemigroup of A (since B-metrics can be considered merely as a particular case of A-metrics with values in B). Suppose our theorem is true for denumerable ordered involuted semigroups. Then every finitely generated ordered involuted subsemigroup B of A satisfies the conditions of the theorem, i.e. $B \in \mathcal{K}$. It is known [10], [11] and [14] that the class \mathcal{K} is universal. Therefore [18], \mathcal{K} has the local property: if every finitely generated subsystem of a system belongs to \mathcal{K} , the system itself belongs to \mathcal{K} . So $A \in \mathcal{K}$. It follows that we need to prove the theorem for the denumerable case only. Therefore, we suppose that A is a denumerable ordered involuted semigroup satisfying the condition of our theorem.

We are going to construct sequences $M_1 < M_2 < \dots$ of nets and $d_1 \subset d_2 \subset \dots$ of proper A-metrics (here d_i is a metric on $M_i, M_{i-1} < M_i$ means that M_{i-1} is a subnet of $M_i, d_{i-1} \subset d_i$ means d_{i-1} is a restriction of d_i) satisfying the following condition (P):

(P) for every chain μ_{ik} in M_m leading from i to k and for every $a_1, a_2 \in A$ such that $d_m(\mu_{ik}) \leq a_1 a_2$ there exist an n and a vertex j in M_n such that $m \leq n$ and $d_n((i,j)) \leq a_1, \ d_n((j,k)) \leq a_2$.

Let N be the set of all natural numbers, and let $(N \times N)^*$ denote the set of all finite sequences of pairs of natural numbers. Then the set $(N \times N)^* \times A \times A$ is infinitely denumerable, so let α be a fixed bijection of this set onto N.

We define M_n , d_n inductively. Choose an element $a \in A$ and define d(0) = a. Then d_1 is a proper A-metric on M_1 .

Suppose that we have defined the sequences $M_1 < ... < M_{n-1}$ of nets and $d_1 \subset ... \subset d_{n-1}$ of metrics. Let $R_{n-1} = \{\alpha(\mu, a_1, a_2): \mu \text{ be a chain in } M_{n-1} \text{ leading, say, from } i \text{ to } k; a_1, a_2 \in A; d_{n-1}(\mu) \leqslant a_1a_2; \text{ there exists no vertex } j \text{ in } M_{n-1} \text{ such that } d_{n-1}((i,j)) \leqslant a_1 \text{ and } d_{n-1}((j,k)) \leqslant a_2\}.$ Notice that any chain in M_{n-1} is a sequence of edges which are ordered pairs of natural numbers (vertices), and so every chain is an element of $(N \times N)^*$ and the definition of R_{n-1} makes sense.

Two cases are possible.

Case 1. R_{n-1} is empty. This means that the sequences $M_1 < \dots < \dots < M_{n-1}$ and $d_1 \subset \dots \subset d_{n-1}$ satisfy condition (P) and we need not prolong the sequences.

Case 2. R_{n-1} is not empty. In this case define $\beta_{n-1} = \min R_{n-1}$. Then $\beta_{n-1} = a(\mu_{ik}, a_1, a_2)$, where μ_{ik} is a chain in M_{n-1} leading from i to k, $d_{n-1}(\mu_{ik}) \leq a_1 a_2$ and there is no vertex j in M_{n-1} such that $d_{n-1}((i,j)) \leq a_1$ and $d_{n-1}((j,k)) \leq a_2$. Define M_n and d_n as follows: add a new vertex n and two new arrows (i,n), (n,k) to M_{n-1} and extend d_{n-1} to d_n supposing $d_n((i,n)) = a_1$, $d_n((n,k)) = a_2$. Clearly, d_n is a proper metric on M_n .

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Proceeding as above, we either obtain finite sequences satisfying property (P) or find that the sequences $M_1 < ...$ and $d_1 \subset ...$ are infinite. However, in the latter (infinite) case property (P) is satisfied. In effect, it follows from the definition of M_n that $\beta_{n-1} < \beta_n$; therefore, $n-1 \leqslant \beta_n$. Suppose $d_m(\mu_{ik}) \leq a_1 a_2$ for a chain μ_{ik} in M_m leading from i to k and for $a_1, a_2 \in A$. Let $a(\mu_{ik}, a_1, a_2) = p$ and n = p + 2. Then, by the minimality of β_n and the evident inequality $\beta_n > p$, there exists a vertex i in M_n such that $d_n((i,j)) \leqslant a_1$ and $d_n((j,k)) \leqslant a_2$.

Let X_a be the set of all vertices of the nets M_i from the sequence constructed above (so, if the sequence is finite, Xu is an initial sequence of N, otherwise, X_a coincides with N). For every $b \in A$ define a binary relation $P_a(b)$ on X_a as follows: $(i,j) \in P_a(b)$ if and only if there exists a net M_n in our sequence and a chain μ_{ij} in M_n leading from i to j such that $d_n(\mu_{ij}) \leq b$.

The following chain of equivalences shows that $P_a(b^{-1}) = P_a(b)^{-1}$:

$$(i,j) \in P_a(b^{-1}) \Leftrightarrow d_n(\mu_{ij}) \leqslant b^{-1} \Leftrightarrow d_n(\mu_{ij})^{-1} \leqslant b \Leftrightarrow d_n(\mu_{ij}^{-1}) \leqslant b \Leftrightarrow (j,i) \in P_a(b) .$$

Here we have used successively the definition of $P_a(b^{-1})$, (6), the definition of the length of a chain and the fact that μ_{ii}^{-1} is a chain in M_n leading from j to i.

Now suppose $(i,j) \in P_a(a_1)$ and $(j,k) \in P_a(a_2)$. Then $d_p(\mu_{ij}) \leq a_1$ and $d_a(\mu_{ik}) \leq a_i$ for chains μ_{ij} in M_p and μ_{ik} in M_q . Let $n = \max\{p, q\}$. Then $d_n(\mu_{ij}) \leqslant a_i$ and $d_n(\mu_{ik}) \leqslant a_i$. Using (5), we obtain: $d_n(\mu_{ij}\mu_{jk}) = d_n(\mu_{ij}) d_n(\mu_{jk})$ $\leq a_1 a_2$ and, since $\mu_{ij}\mu_{jk}$ is a chain in M_n leading from i to k, $(i, k) \in P_a(a_1 a_2)$.

Conversely, suppose $(i, k) \in P_a(a_1 a_2)$. Then $d_m(\mu_{ik}) \leq a_1 a_2$ for some m and a chain μ_{ik} in M_m leading from i to k. By property (P), there exist an n and a vertex j in M_n such that $d_n((i,j)) \leqslant a_1$ and $d_n((j,k)) \leqslant a_2$. Since (i, j) and (j, k) are chains in M_n , $(i, j) \in P_a(a_1)$ and $(j, k) \in P_a(a_2)$. Therefore, $P_a(a_1 a_2) = P_a(a_2) \circ P_a(a_1)$.

If $a_1 \leqslant a_2$ and $(i,j) \in P_a(a_1)$, then $d_n(\mu_{ij}) \leqslant a_1$ for some n and μ_{ij} . It follows that $d_n(\mu_{ij}) \leqslant a_2$ and $(i,j) \in P_a(a_2)$. Thus $a_1 \leqslant a_2$ implies $P_a(a_1)$ $\subset P_a(a_2)$.

Now let X be the sum of the family $(X_a)_{a \in A}$ of sets. Without loss of generality we may suppose that the sets X_a are disjoint (e.g. we may replace X_a by $X_a \times \{a\}$). Define $P(b) = \Sigma(P_a(b))_{a \in A}$ where Σ denotes the disjoint union of $P_a(b)$ (if the sets X_a are disjoint, Σ is the ordinary settheoretical union). Clearly,

$$(9) P(b^{-1}) = P(b)^{-1},$$

(10)
$$P(a_1 a_2) = P(a_2) \circ P(a_1)$$

for all $b, a_1, a_2 \in A$, since the same properties hold for P_a . In the same. way, $a_1 \leqslant a_2 \Rightarrow P(a_1) \subset P(a_2)$. Now suppose $P(a_1) \subset P(a_2)$ for some $a_1, a_2 \in A$.

Then $P_a(a_1) \subseteq P_a(a_2)$ for all $a \in A$. In particular, $P_{a_1}(a_1) \subseteq P_{a_1}(a_2)$. Since $d_1((0,1)) = a_1$, we conclude that $(0,1) \in P_{a_1}(a_1)$; hence $(0,1) \in P_{a_1}(a_2)$. Therefore, there exist an n and a chain μ_{01} in M_n leading from 0 to 1 and such that $d_n(\mu_{01}) \leqslant a_2$. By the condition of our theorem, the metric d_n is regular. Therefore, $a_1 = d_n(\mathbf{0}) \leqslant d_n(\mu_{\text{ot}}) \leqslant a_2$. Thus

$$a_1\leqslant a_2\Leftrightarrow P(a_1)\subset P(a_2).$$

Formulas (9)-(1.1) show that P is an isomorphism of A onto an ordered involuted semigroup of binary relations. In effect, if $P(a_1)$ $=P(a_2)$, then $P(a_1) \subset P(a_2)$ and $P(a_2) \subset P(a_1)$. By (11), $a_1 \leqslant a_2$ and $a_2 \leqslant a_1$; hence $a_1 = a_2$ and P is one-to-one. The form of (10) is due to the fact that we write factors in products from left to right in abstract semigroups (as in A) and from right to left in semigroups of binary relations.

Thus the ordered involuted semigroup is representable, which ends the proof of Theorem 1.

Remark 1. Suppose that $A = (A; \cdot, \cdot^{-1}, \leq)$ is a quasiordered involuted semigroup, i.e. that \leqslant is a *quasiorder* (that is, reflexive and transitive) binary relation on an involuted semigroup $(A; \cdot, \cdot^{-1})$ and properties (5) and (6) are satisfied. A is called representable in a wider sense if there exists a mapping P of A onto an ordered involuted semigroup of binary relations satisfying (9)-(11). This being the case.

$$P(a_1) = P(a_2) \Leftrightarrow (a_1 \leqslant a_2 \& a_2 \leqslant a_1),$$

i.e. P is an isomorphism if and only if \leqslant is an order relation. The same proof as that of Theorem 1 shows that a quasiordered involuted semigroup is representable in a wider sense if and only if it satisfies the condition of Theorem 1 (i.e. every proper metric on every net is regular). An obvious readjustment of the definition of proper metrics in the case of quasiorders is needed.

Remark 2. Let $(A; \cdot, \cdot)$ be an involuted semigroup and let ϱ be a binary relation on A. Define proper A-metrics for $A = (A; \cdot, \cdot^{-1}, \varrho)$ precisely as for ordered involuted semigroups, using ϱ instead of \leqslant . Define $a_1 \leq_o a_2$ for $a_1, a_2 \in A$ if and only if there exist a net M_n and a proper A-metric d on M_n such that $d(0) = a_1$ and $d(\mu) = a_2$ for some chain μ in M_n leading from 0 to 1. Then \leq_o is the smallest element in the set of all quasiorder relations π on A for which $\varrho \subset \pi$ (i.e. $a_1 \varrho \ a_2 \Rightarrow a_1 \pi \ a_2$ for all $a_1, a_2 \in A$) and $(A; \cdot, \cdot^{-1}, \pi)$ is a quasiordered involuted semigroup representable in a wider sense. The proof of this fact is quite analogous to the proofs of Theorems 1 and 2.

Now we are going to characterize representable involuted semigroups. The sense of our necessary and sufficient condition of representability (Theorem 2) is quite simple: on an involuted semigroup A



we construct a quasiorder relation \leq_{-} (cf. Remark 2) which is the smallest quasiorder relation turning A into a quasiordered involuted semigroup representable in a wider sense; A is representable if and only if \leq_{-} is antisymmetric (i.e. \leq_{-} is an order relation). In other words, A is representable if and only if for every two nets M_m and M_n and every two proper A-metrics d_m on M_m and d_n on M_n (here in the definition of a proper metric the equality relation = is used instead of \leq), every two chains μ in M_m and ν in M_n , leading from 0 to 1, the equalities $d_m(\mu) = d_n(0)$ and $d_n(\nu) = d_m(0)$ imply $d_m(0) = d_n(0)$. However, it turns out that we may assume $M_m = M_n$ in the above condition. Moreover, if every element of A is decomposable into a product of two elements, the metrics d_m and d_n in the above condition may be chosen alike (in the sense that the chains figuring in the definition of d_m are the same as the chains for d_n). To make all elements of A decomposable we add a new element to A.

Let $A=(A;\cdot,^{-1})$ be an involuted semigroup. If the semigroup $(A;\cdot)$ contains an identity 1, then $1^{-1}=1$. In effect, $1^{-1}=1\cdot 1^{-1}=(1\cdot 1^{-1})^{-1}=(1^{-1})^{-1}=1$. If $(A;\cdot)$ does not contain an identity, adjoin a new element 1 to A and define $1\cdot x=x\cdot 1=x$ for all $x\in A\cup\{1\}, 1^{-1}=1$. As a result we obtain a new involuted semigroup with identity; this involuted semigroup will be denoted by A^1 . If A contains an identity, suppose $A^1=A$.

Suppose that A is an involuted semigroup, and M_n is a net. An A-metric on M_n is any mapping d of the set $\{0,1,...,2n-2\}$ of all arrows of M_n into A. Lengths of chains in any A-metric d on M_n are defined precisely as in the case of ordered involuted semigroups. An A-metric d on M_n is called proper if for any k, $0 < k \le n$, there exists a chain μ_k in M_k leading from p_{k+1} to q_{k+1} whose length equals that of the chain (2k-1,2k), i.e. $d(\mu_k)=d(2k-1)d(2k)$. Two proper A-metrics d and d' on M_n are called similar if for any k, 0 < k < n, there exists a chain μ_k in M_k leading from p_{k+1} to q_{k+1} and such that $d(\mu_k)=d(2k-1)d(2k)$, $d'(\mu_k)=d'(2k-1)d'(2k)$. In other words, two proper A-metrics are similar if the same chains can be used in ascertaining the properness of both metrics.

If d is a metric on A, then d(0) is called the major value of d. Two metrics d and d' on M_n are called chained if there exists a chain μ in M_n leading from 0 to 1 and such that $d(\mu) = d'(0)$, $d'(\mu) = d(0)$.

THEOREM 2. An involuted semigroup A is representable if and only if every two chained similar proper A^1 -metrics on every net have the same major values.

Proof. Necessity. As in the proof of Theorem 1, it suffices to prove the condition of our theorem for an involuted semigroup F of binary relations. We shall prove a stronger condition: for every two nets M_m

and M_n and chains μ in M_m , ν in M_n , both leading from 0 to 1, and for every two proper F^1 -metrics, d_m on M_m and d_n on M_n , the equalities $d_m(\mu) = d_n(0)$ and $d_n(\nu) = d_m(0)$ imply $d_m(0) = d_n(0)$.

Suppose that M_m , M_n , μ , ν , d_m and d_n are as above and the above two equalities hold. Since F^1 is an involuted semigroup of binary relations (*), it is ordered by the inclusion order relation C; let F^C denote the corresponding ordered involuted semigroup. Now d_m and d_n may be considered as F^C -metrics; obviously, both d_m and d_n are proper F^C -metrics. By Theorem 1, $d_m(0) \subset d_m(\mu)$ and $d_n(0) \subset d_n(\nu)$; hence $d_m(0) \subset d_m(\mu) = d_n(0)$ and analogously $d_n(0) \subset d_m(0)$. It follows that $d_m(0) = d_n(0)$.

Sufficiency. One possible approach is to define a binary relation \leq on the set A of elements of an involuted semigroup A: $a_1 \leq a_2$ if and only if there exist a net M_n , a proper A^1 -metric d on M_n and a chain μ in M_n leading from 0 to 1 and such that $a_1 = d(0)$, $a_2 = d(\mu)$. Then we could prove that $(A; \cdot, \neg^1, \leq)$ was an ordered involuted semigroup satisfying the condition of Theorem 1. Thus $(A; \cdot, \neg^1, \leq)$ would be representable; hence A would be representable as well. However, such an approach leads to an entirely new proof as compared with that of Theorem 1. To save room we give another proof, using that of Theorem 1.

As in Theorem 1, we see that without loss of generality A may be supposed to be a finite or countably infinite involuted semigroup satisfying the condition of Theorem 2. Now copy the whole proof of Theorem 1 up to the formulas (9) and (10) with the following alterations: consider A^1 instead of A and = instead of \leq . The proof is transferred to this case without alterations. Then P is a homomorphism of A^1 onto an involuted semigroup of binary relations. It remains to prove that P is an isomorphism, i.e. that P is one-to-one.

Suppose $P(a_1) = P(a_2)$ for $a_1, a_2 \in A$. Then $P_a(a_1) = P_a(a_2)$ for all $a \in A$. In particular, $P_{a_1}(a_1) = P_{a_1}(a_2)$. As in the proof of sufficiency of Theorem 1, we obtain $(0, 1) \in P_{a_1}(a_2)$, i.e. there exist an n and a chain μ in M_n leading from 0 to 1 such that $d_n(0) = a_1$ and $d_n(\mu) = a_2$ for a proper A^1 -metric d_n on M_n . Analogously, there exist an m, a chain ν in M'_m and a proper A^1 -metric d'_m on M'_m such that $d'_m(0) = a_2$ and $d'_m(\nu) = a_1$. Here condition (P) is satisfied for M_n , d_n and for M'_m , d'_m .

We are going to construct a new net M of order m+n "glueing together" M_n and M'_m as in Fig. 2, where the arrows displayed in M_n and M'_m are the 0th arrows and a new arrow is added. The problem is to renumber the vertices and arrows of M.

The vertices of M are $\{0, 1, ..., m+n\}$, the arrows are $\{0, 1, ..., 2m++2n-2\}$. Define injections f and g of the sets of vertices of M_n and M'_m into the set of vertices of M as follows: f(0) = 0, f(k) = k+1 for $1 \le k \le n$;

^(*) We suppose $F^1 = F \cup \{\Delta\}$ where Δ is the identity relation in the case of $F^1 \neq F$.



g(0) = 2, g(1) = 1, g(k) = k + n for $2 \le k \le m$. If (i,j) is an arrow in M_n , then (f(i), f(j)) is an arrow in M, if (i,j) is an arrow in M'_m , then (g(i), g(j)) is an arrow in M. In other words, f and g are isomorphic embeddings of M_n and M'_m into M. If (i,j) is an arrow in M_n , then we will write f((i,j)) instead of (f(i), f(j)); g((i,j)) has an analogous meaning. $f((i,j))^{-1} = (f(i), f(j))^{-1}$; $g((i,j))^{-1} = (g(i), g(j))^{-1}$. If π is a chain in M_n (in M'_m), then $f(\pi)$ $(g(\pi))$ denotes the corresponding chain in M.

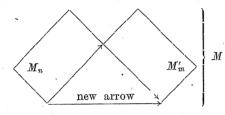


Fig. 2

Define A^1 -metrics d and d' on M as follows: $d(0) = a_1$, $d(f(i)) = d_n(i)$ for every arrow i of M_n , d(g(i)) = 1 for every arrow i of M'_m ; $d'(0) = a_2$, d'(f(i)) = 1 for every arrow i of M_n , $d'(g(i)) = d'_m(i)$ for every arrow i of M'_m . To check the fact that both d and d' are proper we need only mention that $d(1)d(2) = a_1 \cdot 1 = a_1 = d(0)$, $d'(1)d'(2) = 1 \cdot a_2 = a_2 = d'(0)$. Clearly, d and d' are similar, since all the chains we used to show that d_n and d'_m are proper are isomorphically mapped (by means of f or g) into M. Since $d(f(\mu), g(v)) = d(f(\mu))d(g(v)) = a_2 \cdot 1 = a_2 = d'(0)$ and $d'(f(\mu), g(v)) = d'(f(\mu))d'(g(v)) = 1 \cdot a_1 = a_1 = d(0)$, the metrics d and d' are chained. Here $(f(\mu), g(v))$ denotes the juxtaposing of the chains $f(\mu)$ and g(v) in M. Clearly, it is a chain in M leading from 0 to 1.

By the condition of our theorem, $a_1 = d(0) = d'(0) = a_2$, i.e. P is one-to-one. P induces an isomorphism of A onto an involuted semigroup of binary relations, i.e. A is representable, which completes the proof of Theorem 2.

Suppose that G is a set, $g \mapsto g^*$ is a bijection of G onto a set G^* , G and G^* are disjoint, and $H = G \cup G^*$. The set FI(G) of all nonempty words over the alphabet H is an involuted semigroup under the juxtaposition of words and the following involution: $(h_1h_2 \dots h_n)^{-1} = h_n^{-1} \dots h_2^{-1}h_1^{-1}$ for all $h_1, h_2, \dots, h_n \in H$. Here $g^{-1} = g^*$ and $(g^*)^{-1} = g$ for all $g \in G$.

FI(G) is a free involuted semigroup, and G is a set of free generators of FI(G). We identify the identity of $FI(G)^1$ with the empty word 1.

Proposition 3. Free involuted semigroups are representable.

Proof. Suppose that FI(G) is a free involuted semigroup, and |a| denotes the length of the word $a \in FI(G)^1$; in particular, |1| = 0.

Let M_n be a net, and let d and d' be chained similar proper $FI(G)^1$ -metrics on M_n , $d(\mu) = d'(0)$, $d'(\mu) = d(0)$ for a chain μ in M_n leading from 0 to 1.

We shall prove that $|d(\mathbf{0})| \leqslant |d(\mathbf{v})|$ for every chain \mathbf{v} leading from $\mathbf{0}$ to 1, and |d(0)| = |d(v)| implies d(0) = d(v). Let m be the minimal number such that ν is a chain in M_m . Then $m \leqslant n$. We give a proof by induction. If m=1, then ν is of the form $(0,0^{-1},0,...,\tilde{0}^{-1},0)$ and our statement is obviously true. Suppose it is true for all $m \leqslant k$. Let ν be a chain in M_{k+1} but not in M_k . Then at least one of the edges 2k-1, $(2k-1)^{-1}$, 2k, $(2k)^{-1}$ belongs to v. Suppose that μ_k is a chain in M_k leading from p_{k+1} to q_{k+1} and such that $d(\mu_k) = d(2k-1) \cdot d(2k)$. Such a chain does exist since d is proper. Replace every subchain (2k-1, 2k) occurring in ν by μ_k and every subchain $((2k)^{-1}, (2k-1)^{-1}) = (2k-1, 2k)^{-1}$ by μ_k^{-1} . We obtain a chain $\bar{\nu}$ leading from 0 to 1. Clearly, $d(v) = d(\bar{v})$. Omit all the occurrences of subchains $(2k-1, (2k-1)^{-1}), ((2k)^{-1}, 2k)$ from $\bar{\nu}$. We obtain a chain $\bar{\bar{\nu}}$ leading from 0 to 1 and such that $|d(\bar{v})| \leq |d(\bar{v})|$. Clearly, $|d(\bar{v})| = |d(\bar{v})|$ if and only if $\bar{\nu}$ does not contain the above subchains, i.e. $\bar{\nu} = \bar{\nu}$. Now, $\bar{\bar{\nu}}$ is a chain in M_k and, by the supposition of induction, $|d(\mathbf{0})| \leqslant |d(\bar{\bar{\nu}})|$. Therefore, $|d(0)| \leq |d(v)|$. If |d(0)| = |d(v)|, then $|d(0)| = |d(\overline{v})|$, $v = \overline{v}$ and, by the supposition of induction, $d(0) = d(\bar{\nu}) = d(\bar{\nu}) = d(\nu)$.

Thus $|d(0)| \le |d(\mu)| = |d'(0)| \le |d'(\mu)| = |d(0)|$; hence, $|d(0)| = |d(\mu)|$ and, by the above statement, $d(0) = d(\mu) = d'(0)$. Therefore, the condition of Theorem 2 is fulfilled. It follows that Fl(G) is representable.

COROLLARY 4. The class of all representable involuted semigroups does not form a variety of algebras, since it is not closed under homomorphisms.

Proof. Every involuted semigroup is a homomorphic image of a free one; there exist involuted semigroups which are not representable [14]; free involuted semigroups are representable by Proposition 3.

Remark 3. It was proved in [14] and in [10] and [11] that the class of all representable involuted semigroups forms a quasivariety of algebras (i.e. a class characterized by a system of quasiidentities, a quasiidentity being a universal formula of the form $A_1 \& A_2 \& ... \& A_n \Rightarrow A$, where $A_1, A_2, ..., A_n$, A are equations and no quantifiers occur). It follows from Corollary 4 that the class of all representable involuted semigroups cannot be characterized by a system of identities. Moreover, it follows from Proposition 3 that any identity that holds for all representable involuted semigroups holds for all involuted semigroups as well.

The conditions of representability given in Theorems 1 and 2 do not have the form of conditional identities. However, they can easily be transformed into conditional identities (i.e. quasiidentities), which we do in what follows.

Let M_n be a net. Consider the first order predicate calculus fit for



(ordered) involuted semigroups. The alphabet of the calculus comprises a countably infinite list of individual variables x_0, x_1, \ldots , two symbols for the operations of involuted semigroups (and a symbol for the order relation), logical symbols (propositional connectives, quantifiers), equality, technical symbols (brackets). With every arrow i of M_n associate the individual variable x_i . With every chain $\mu = (e_1, e_2, \ldots, e_m)$ in M_n associate the term (polynomial) $[\mu]$ as follows: $[i] = x_i, [i^{-1}] = x_i^{-1}, [\mu] = [e_1][e_2] \ldots [e_m]$. For example, $[(0, 4^{-1}, 6^{-1}, 6, 3^{-1}, 1)] = x_0 x_4^{-1} x_0^{-1} x_0 x_3^{-1} x_1$.

A chain system in M_n is a sequence $(\mu_1, \mu_2, ..., \mu_{n-1}, \mu)$ where μ_i is a chain in M_i leading from p_{i+1} to q_{i+1} , μ is a chain in M_n leading from 0 to 1.

With every chain system S in M_n associate the following formula $\langle S \rangle$: if $S = (\mu_1, \mu_2, \dots, \mu_{n-1}, \mu)$, then $\langle S \rangle$ is

$$[\mu_1] \leqslant x_1 x_2 \& [\mu_2] \leqslant x_3 x_4 \& \dots \& [\mu_i] \leqslant x_{2i-1} x_{2i} \& \dots$$

... &
$$[\mu_{n-1}] \leq x_{2n-3} x_{2n-2} \Rightarrow x_0 \leq [\mu]$$
.

THEOREM 5. An ordered involuted semigroup is representable if and only if it satisfies conditions $\langle S \rangle$ for every chain system S in every net M_n .

Proof. Suppose that S is a chain system in a net M_n . Substitute an element d(i) of an ordered involuted semigroup A for each occurrence of x_i in S. Then d can be considered as a metric on M_n . If the antecedent of the implication $\langle S \rangle$ is valid, this means that d is a proper A-metric (for the chains μ_1, \ldots, μ_{n-1} occurring in S). The regularity of d means that $d(0) \leq d(\mu)$ for every chain μ in M_n leading from 0 to 1. In other words, $x_0 \leq \lfloor \mu \rfloor$ for our assignment of values in A for individual variables x_i . Thus A satisfies the condition of Theorem 5 if and only if every proper A-metric on every net M_n is regular. Now Theorem 5 follows from Theorem 1, which completes the proof.

Theorem 5 characterizes the class of all representable ordered involuted semigroups in terms of elementary axioms. Clearly, the system of axioms given in Theorem 5 is infinite. Moreover, for every net M_n we can construct an infinite system of axioms, since there exists an infinite set of chain systems in M_n . The axioms obtained in this way are not independent. Roughly speaking, shorter axioms follow from longer ones. In what follows we reduce the number of axioms corresponding to a given net M_n .

A chain ν in M_n is called *straightened* if it does not contain subchains of the form (2i-1, 2i) or $(2i-1, 2i)^{-1}$.

RESTRICTION 1. In conditions $\langle S \rangle$ associated with chain systems we may consider straightened chains only.

Motivation. Suppose μ_k is a chain in a chain system S and μ_k contains a subchain (2i-1,2i). Clearly, i < k. The chain μ_i leads from

 p_{i+1} to q_{i+1} as well as the chain (2i-1, 2i). Replace (2i-1, 2i) by μ_i in μ_k . We obtain a new chain $\overline{\mu_k}$. Now if the antecedent of $\langle S \rangle$ holds for some $x_i \in A$, then $[\mu_i] \leqslant x_{2i-1}x_{2i}$ and $[\mu_k] \leqslant x_{2k-1}x_{2k}$. By (5), $[\overline{\mu_k}] \leqslant x_{2k-1}x_{2k}$, since $[\overline{\mu_k}] \leqslant [\mu_k]$. Thus μ_k may be replaced by $\overline{\mu_k}$ in the antecedent of $\langle S \rangle$. Let i be the smallest number such that the chain μ_i is not straightened. Clearly, μ_1 is always straightened, and so $2 \leqslant i$. As we have just seen, we may replace μ_i (by replacing all occurrences of subchains of the form (2j-1, 2j) and $(2j-1, 2j)^{-1}$ by μ_j and μ_j^{-1}) by another chain $\overline{\mu_i}$ which is straightened. We obtain a new chain system \overline{S} and condition $\langle \overline{S} \rangle$ implies $\langle S \rangle$. It is now easily seen that we can replace all non-straightened chains in the antecedent of $\langle S \rangle$ by straightened ones, and the resulting condition implies $\langle S \rangle$.

It remains to consider the final chain μ of \mathcal{S} . If μ contains a subchain (2i-1, 2i) and we replace it by μ_i , we obtain a new chain $\overline{\mu}$. If $x_0 \leq \lceil \overline{\mu} \rceil$, then $\lceil \overline{\mu} \rceil \leq \lceil \mu \rceil$ and $x_0 \leq \lceil \mu \rceil$. Therefore, if μ is replaced by $\overline{\mu}$ in $\langle \mathcal{S} \rangle$, we obtain a stronger condition. It follows that we may consider the chain systems with straightened chains only.

Consider the net M_1 . Every chain in M_1 is of the form $(0, 0^{-1}, 0, ...$..., $0^{-1}, 0)$ and every chain system in M_1 is of the form (μ) . Therefore, we obtain the following axioms corresponding to M_1 : $x_0 \leq x_0 x_0^{-1} x_0 ... x_0^{-1} x_0$. In particular,

$$(12) x \leqslant xx^{-1}x$$

(here x_0 is replaced by x for convenience). Clearly, (12) implies every other axiom corresponding to M_1 . By Theorem 5, every representable ordered involuted semigroup satisfies (12).

Let ν be a chain in M_n . Then (ν, ν^{-1}, ν) is a chain which is called a zigzag. If a chain μ does not contain subchains which are zigzags, μ is called a chain without zigzags.

RESTRICTION 2. In the case of ordered involuted semigroups satisfying condition (12) we may consider axioms $\langle S \rangle$ for those chain systems S only which consist of chains without zigzags.

Motivation. Suppose a chain π contains a zigzag (v, v^{-1}, v) . Replacing the zigzag by v and doing the same with all other zigzags, we obtain a chain π without zigzags. Clearly, $[\pi] \leq [\pi]$ is true in every ordered involuted semigroup satisfying (12). It follows that if we replace all chains π by π in S, we obtain a new chain system \overline{S} such that $\langle \overline{S} \rangle$ implies $\langle S \rangle$.

It can be proved that for every net M_n there exists only a finite number of different chain systems satisfying restrictions 1 and 2. Thus every net produces a finite number of axioms for the class of representable ordered involuted semigroups.



It can also be proved that the only new axiom corresponding to M_2 is $x_0 \leqslant x_1 x_2 \Rightarrow x_0 \leqslant x_1 x_1^{-1} x_0$ (all the other axioms follow from this one and from (5), (6) and (12)). The only new axiom which we obtain considering various nets M_3 is $x_0 \leqslant x_1 x_2 \& x_0 x_2^{-1} \leqslant x_3 x_4 \Rightarrow x_0 \leqslant x_1 x_4^{-1} x_4 x_2$. The other axioms follow from (12) and the ones just mentioned.

Now we will give a system of quasiidentities for the class of representable involuted semigroups. To this end we suppose that the language of the lower predicate calculus which we use contains the individual variables x_0, x_1, \ldots and y_0, y_1, \ldots ; if μ is a chain in M_n , then $[\mu]_x$ denotes the polynomial constructed above, while $[\mu]_y$ denotes the polynomial constructed in the above fashion where all variables x_i are replaced by y_i . If $S = (\mu_1, \ldots, \mu_{n-1}, \mu)$ is a chain system in a net M_n , then (S) is the following quasiidentity:

$$\begin{split} & [\mu_1]_x = x_1 x_2 \ \& \dots \ \& \ [\mu_{n-1}]_x = x_{2n-3} x_{2n-2} \ \& \ [\mu]_x = y_0 \ \& \\ & \& \ [\mu_1]_y = y_1 y_2 \ \& \dots \ \& \ [\mu_{n-1}]_y = y_{2n-3} y_{2n-2} \ \& \ [\mu]_y = x_0 \Rightarrow x_0 = y_0 \ . \end{split}$$

As in the proof of Theorem 5, we can show that conditions (S) for all chain systems S in all nets are equivalent to the conditions of Theorem 2; hence, the following result holds:

THEOREM 6. An involuted semigroup A is representable if and only if A^1 satisfies the quasiidentities (S) for every chain system S in every net M_n .

As in the case of ordered involuted semigroups, we may consider quasiidentities corresponding to straightened chain systems only.

We have characterized the classes of representable ordered and unordered involuted semigroups by infinite systems of elementary axioms. We conjecture that the classes are not finitely axiomatizable. It would be interesting to find the axioms for our classes of algebraic systems in the form of explicit axiom schemata without any use of "graph-theoretic scaffolding" consisting of nets.

The characterization of representable (ordered) semigrouds which we give below is quite similar to that of (ordered) involuted semigroups; therefore we will restrict ourselves to outlines of the proofs.

First of all we should modify appropriately the concept of net. A net of order n $(0 \le n)$ is a graph B_n with the vertices $\{0, 1, ..., 2n+1\}$ and arrows $\{0, 1, ..., 3n\}$; each arrow leads from an even vertex to an odd one; (0, 1) is an arrow; for every $i, 1 \le i \le n$, there exist uniquely defined vertices $p_i, q_i \in \{0, 1, ..., 2i-1\}$ such that $(p_i, 2i+1), (2i, 2i+1), (2i, q_i)$ are arrows. It follows that p_i is even and q_i is odd. The vertices $\{0, 1, ..., 2i+1\}$ and the arrows $\{0, 1, ..., 3i\}$ form a subgraph which is a net. It is called a subnet (of order i). We denote this fact as follows:

 $B_i < B_n$. Thus there exists a chain of subnets $B_0 < B_1 < ... < B_n$. Clearly, a net is a bipartite graph.

Thus B_n is constructed from B_{n-1} by adding two new vertices and three new arrows, as is shown in Fig. 3

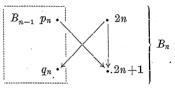


Fig. 3

In B_n we will consider only those chains which lead from an even vertex to an odd one. It is easily seen that every such chain consists of an odd number of edges.

If $A=(A; [\], \leqslant)$ is an ordered semigroud, an A-metric on a net B_n is any mapping $d\colon \{0,1,\ldots,3n\}\to A$. An A-metric d is called proper if for every $i,1\leqslant i\leqslant n$, there exists a chain μ_i in B_{i-1} leading from p_i to q_i and such that $d(\mu_i)\leqslant [d(3i-2)d((3i-1)^{-1})d(3i)]$. Here $3i-2=(p_i,2i+1),\ 3i-1=(2i,2i+1),\ 3i=(2i,q_i)$ and $d(\mu_i)$ is defined in the obvious fashion: if $\mu_i=(e_1,e_2,\ldots,e_{2k-1})$ then $d(\mu_i)=[d(e_1)d(e_2^{-1})\ldots d(e_{2k-2}^{-1})d(e_{2k-1})]$. Here the edges with odd indices are arrows and edges with even indices are reversed arrows, the product of 2k-1 factors is defined inductively: if k>1, then $[a_1\ldots a_{2k-1}]=[[a_1\ldots a_{2k-3}]a_{2k-2}a_{2k-1}]$.

A proper A-metric d on B_n is called regular if $d(0) \leq d(\mu)$ for every chain μ in B_n leading from 0 to 1. Here 0 = (0, 1).

THEOREM 7. An ordered semigroud A is representable if and only if every proper A-metric on every net is regular.

Proof. Suppose $F = (F; [\], \subset)$ is an ordered semigroud of binary relations between the elements of two sets X and Y. Let d be a proper F-metric on a net B_n and $(x_0, x_1) \in d(0)$ for some $x_0 \in X$ and $x_1 \in Y$. We are going to construct a sequence $x_0, x_1, \ldots, x_{2n+1}$ such that the terms with even indices belong to X, the terms with odd indices belong to Y and $(x_{2i}, x_{2i+1}) \in d((2i, 2i+1))$ for all $i, 0 \leq i \leq n$. We are given x_0 and x_1 . Suppose that the elements $x_0, x_1, \ldots, x_{2i-1}$ satisfying the above condition have already been chosen. Then

$$(p_i, q_i) \in d(\mu_i) \subset [d(3i-2)d((3i-1)^{-1})d(3i)],$$

i.e. there exist $x_{2i+1} \in Y$ and $x_{2i} \in X$ such that $(p_i, 2i+1) \in d(3i-2)$, $(2i, 2i+1) \in d(3i-1)$, $(2i, q_i) \in d(3i)$. Proceeding as above, we choose the terms of the sequence x_0, \ldots, x_{2n+1} .

If μ is a chain in B_n leading from 0 to 1, then $(x_0, x_1) \in d(\mu)$. Since this is true for arbitrary $x_0 \in X$ and $x_1 \in Y$, $d(0) \subset d(\mu)$, i.e. d is regular.

Sufficiency. As in the proof of Theorem 1, we see that it suffices to prove the representability of a finite or countably infinite ordered semigroud A satisfying the condition of Theorem 7. We construct a finite or infinite sequence $B_0 < B_1 < ...$ of nets and $d_0 \subset d_1 \subset ...$ of proper A-metrics $(d_i$ is a metric on B_i) such that the following condition holds:

(Q) for every chain μ in B_m leading from 2i to 2j-1 and for every $a_1, a_2, a_3 \in A$ such that $d(\mu) \leq [a_1 a_2 a_3]$ there exist an n and vertices 2p, 2q-1 in B_n such that $m \leq n$ and $d_n((2i, 2q-1)) \leq a_1, \ d_n((2p, 2q-1)) \leq a_2, \ d_n((2p, 2j-1)) \leq a_3.$

Such sequences may be constructed in a way analogous to that used for constructing sequences satisfying property (P) in the proof of sufficiency of Theorem 1.

Suppose that such sequences have been constructed and $d_0(0) = a$. For every $b \in A$ define a binary relation $P_a(b)$ between even and odd numbers: $(2i, 2j-1) \in P_a(b)$ iff $d_n(\mu) \leq b$ for some n, a chain μ in B_n leading from 2i to 2j-1. One can easily verify that

$$P_a([a_1 a_2 a_3]) = [P_a(a_1) P_a(a_2) P_a(a_3)] \text{ and } a_1 \leq a_2 \Rightarrow P_a(a_1) \subset P_a(a_2).$$

Suppose that P(b) is a disjoint union (sum) of binary relations $P_a(b)$ for all $a \in A$. Then $P([a_1 a_2 a_3]) = [P(a_1) P(a_2) P(a_3)]$ and $a_1 \leqslant a_2 \Rightarrow P(a_1) \subset P(a_2)$. Now if $P(a_1) \subset P(a_2)$, then $(0,1) \in P_{a_1}(a_1) \subset P_{a_1}(a_2)$; therefore, $d_n(\mu) \leqslant a_2$ for some chain in B_n leading from 0 to 1. By the regularity of d_n , $d_n(0) \leqslant d_n(\mu) \leqslant a_2$; therefore, $a_1 = d_n(0) \leqslant a_2$, i.e. $a_1 \leqslant a_2 \Leftrightarrow P(a_1) \subset P(a_2)$. It follows that P is an isomorphism of A onto an ordered semigroud of binary relations, i.e. A is representable. Theorem 7 is proved.

A chain system in a net B_n is a sequence $(\mu_1, ..., \mu_n, \mu)$ where μ_i is a chain in B_{i-1} leading from p_i to q_i and μ is a chain in B_n leading from 0 to 1. To every chain $v = (e_1, e_2, ..., e_{2k-1})$ there corresponds a polynomial $[v] = [[e_1][e_2] ... [e_{2k-1}]]$ where $[e_i] = x_j$ if $e_i = j$ or $e_i = j^{-1}$. If S is a chain system in B_n , then $\langle S \rangle$ denotes the following formula:

$$[\mu_1] \leqslant [x_1 x_2 x_3] \& [\mu_2] \leqslant [x_4 x_5 x_6] \& \dots \& [\mu_n] \leqslant [x_{9n-2} x_{9n-1} x_{9n}] \Rightarrow x_0 \leqslant [\mu].$$

THEOREM 8. An ordered semigroud is representable if and only if it satisfies the axioms $\langle S \rangle$ for all chain systems S in all nets B_n .

This theorem can be deduced from Theorem 7 in the same way as Theorem 5 has been deduced from Theorem 1.

In what follows $[\mu]_x$ denote $[\mu]$ in the above sense and $[\mu]_y$ denotes an analogous polynomial in which x_i are replaced by y_i .

THEOREM 9. A semigroud is representable if and only if for all chain

systems $S = (\mu_1, ..., \mu_m, \mu)$ and $T = (\nu_1, ..., \nu_n, \nu)$ in arbitrary nets B_m and B_n respectively the following condition (S, T) is satisfied:

$$\begin{aligned} & [\mu_1]_x = [x_1x_2x_3] \;\& \ldots \;\& \; [\mu_m]_x = [x_{3m-2}x_{3m-1}x_{3m}] \;\& \; [\mu]_x = y_0 \;\& \\ & \& \; [r_1]_y = [y_1y_2y_3] \;\& \ldots \;\& \; [r_n]_y = [y_{3n-2}y_{3n-1}y_{3n}] \;\& \; [r]_y = x_0 \Rightarrow x_0 = y_0 \;. \end{aligned}$$

Proof. It can easily be verified that the condition of Theorem 9 means that for all nets B_m and B_n , all proper A-metrics d_m (for a semi-groud A) and d_n on B_m and B_n , respectively, and all chains μ and ν in B_m and B_n , each leading from 0 to 1, the following condition is satisfied: if $d_m(\mu) = d_n(0)$ and $d_n(\nu) = d_m(0)$, then $d_m(0) = d_n(0)$. The proof that a semigroud A is representable if and only if it satisfies the latter condition is quite analogous to the proof of Theorem 2 and will be omitted for this reason.

Remark 4. Theorem 9 is analogous to Theorem 6; however, conditions (S,T) of Theorem 9 are more complicated than those of Theorem 6, since two nets and two different chain systems S and T are involved. In Theorems 2 and 6 we circumvented this complication by adding an identity to A and considering A^1 -metrics (in particular, the axioms of Theorem 6 should be satisfied by A^1). We could do the same in Theorem 9, considering the conditions of the form (S,S); however, in this case they should be satisfied not in A but in an appropriate extension of A (a new element u such that [uux] = [xuu] = x for all $x \in A \cup \{u\}$ should be added, unless A already has one; this implies that an infinity of new elements of the form $[ux_1ux_2u \dots ux_nu]$ should be added).

A chain μ is called straightened if it does not contain subchains of the form $(3i-2,(3i-1)^{-1},3i)$ or $(3i-2,(3i-1)^{-1},3i)^{-1}$; μ is called without zigzags if it does not contain subchains of the form (v,v^{-1},v) . One can easily verify (using B_0 only) that a representable ordered semigroud satisfies the condition

$$(13) x \leqslant [xxx].$$

As in Restrictions 1 and 2, it can be proved that in Theorems 8 and 9 we may consider straightened chains only; in Theorem 8 we may consider chains without zigzags only, provided the ordered semigroud satisfies (13). The following analogue of Proposition 3 and Corollary 4 holds:

PROPOSITION 10. Free semigrouds are representable. The class of all representable semigrouds is not a variety of algebras since it is not closed under homomorphisms.

It can be seen that a free semigroud with a set X of free generators is merely the set of all words of odd length over the alphabet X with the following ternary operation:

$$[(x_1 \dots x_{2i-1})(y_1 \dots y_{2j-1})(z_1 \dots z_{2k-1})] = x_1 \dots x_{2i-1} y_{2j-1} \dots y_1 z_1 \dots z_{2k-1} \cdot \dots \cdot y_1 z_1 \dots z_{2k-1} \cdot \dots z_{2k-1} \cdot$$



It has been noted in the introduction to this paper that to every involuted semigroup $A=(A; \cdot, \cdot^{-1})$ there corresponds an associated semigroud $[A]=(A; [\cdot])$ where $[a_1a_2a_3]=a_1a_2^{-1}a_3$. In the same way, to every ordered involuted semigroup A there corresponds an associated ordered semigroud [A] with the same order relation as in A.

Proposition 11. An (ordered) involuted semigroup with identity is representable if and only if the associated (ordered) semigroud is representable.

Proof. The "if" part follows from the obvious fact: isomorphisms of (ordered) involuted semigroups are isomorphisms of the associated (ordered) semigrouds as well.

To prove the "if" part suppose that A is an (ordered) involuted semigroup with the identity 1 and [A] is representable. Let P be an isomorphism of [A] onto an (ordered) semigroud of binary relations between the elements of two sets X and Y. Define $Q(a) = P(1)^{-1} \circ P(a)$ for all $a \in A$. Then Q(a) is a binary relation on X. Now $Q(ab) = Q(a1b) = Q([a1b]) = P(1)^{-1} \circ P([a1b]) = P(1)^{-1} \circ P(b) \circ P(1)^{-1} \circ P(a) = Q(b) \circ Q(a); \quad Q(a^{-1}) = Q(1a^{-1}1) = Q([1a1]) = P(1)^{-1} \circ P([1a1]) = P(1)^{-1} \circ P(1) \circ P(a)^{-1} \circ P(1) = P(a) \circ P(1^{-1} \circ P(a))^{-1} \circ P(1) = P([11a])^{-1} \circ P(1) = P(a)^{-1} \circ P(1) = P(a)^{-1} \circ P(a) = P(1) \circ P(a)^{-1} \circ P(a) = P(1) \circ Q(a) = P(1) \circ Q(b) = P(1) \circ P(1)^{-1} \circ P(b) = P([b11]) = P(b), \text{ hence, } a = b. \text{ If } A \text{ is ordered, then } a \leqslant b \Rightarrow P(a) \subset P(b) \Rightarrow P(1)^{-1} \circ P(a) \subset P(1)^{-1} \circ P(b) \Rightarrow Q(a) \subset Q(b) \Rightarrow P(1) \circ Q(a) \subset P(1) \circ Q(b) \Rightarrow P([a11]) \subset P([b11]) \Rightarrow P(a) \subset P(b) \Rightarrow a \leqslant b, \text{ i.e. } a \leqslant b \Rightarrow Q(a) \subset Q(b). \text{ It follows that } Q \text{ is an isomorphism of } A \text{ onto an (ordered) semigroud of binary relations. Proposition 11 is proved.}$

It would be interesting to know whether Proposition 11 is valid for arbitrary (ordered) involuted semigroups.

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