

## Note on a theorem of J. Baumgartner

by

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J. Baumgartner has proved that if  $V = L$  (the Axiom of Constructibility) is assumed, then there is an Aronszajn tree order embeddable in the reals but not in the rationals. We extend this result to show that, under a weaker assumption than  $V = L$ , there are  $2^{\aleph_1}$  non-isomorphic such trees.

We wish to thank Richard Laver for introducing us to the topic here discussed.

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**1. Introduction.** We work in Zermelo-Fraenkel set theory (including Choice), denoted by ZFC, and use the usual notations and conventions.  $R$  denotes the real numbers (as an ordered set) and  $Q$  denotes the rationals.

A *tree* is a poset  $T = \langle T, \leq_T \rangle$  such that for any  $x \in T$ ,  $\text{pr}(x) = \{y \in T \mid y <_T x\}$  is well-ordered by  $\leq_T$ . The order-type of  $\text{pr}(x)$  is the *height* of  $x$  in  $T$ ,  $\text{ht}(x)$ . For each ordinal  $\alpha$  we set  $T_\alpha = \{x \in T \mid \text{ht}(x) = \alpha\}$  and  $T \upharpoonright \alpha = \bigcup_{\beta < \alpha} T_\beta$ .  $T_\alpha$  is the  $\alpha$ th *level* of  $T$ . A *branch* of  $T$  is a maximal totally ordered subset of  $T$ ; if it has order-type  $\alpha$  it is an  $\alpha$ -*branch*. An *antichain* of  $T$  is a pairwise incomparable subset of  $T$ .

Let  $\lambda \leq \omega_1$ . A tree  $T$  is a  $\lambda$ -*tree* if:

- (i)  $(\forall \alpha < \lambda)(T_\alpha \neq \emptyset) \ \& \ T_\lambda = \emptyset$ ;
- (ii)  $(\forall \alpha < \lambda)(|T_\alpha| \leq \aleph_0) \ \& \ |T_0| = 1$ ;
- (iii)  $(\forall \alpha < \lambda)(\forall x \in T_\alpha)(|\{y \in T_{\alpha+1} \mid x <_T y\}| = \aleph_0) \vee (\alpha + 1 = \lambda)$ ;
- (iv)  $(\forall \alpha < \beta < \lambda)(\forall x \in T_\alpha)(\exists y \in T_\beta)(x <_T y)$ ;
- (v)  $(\forall \alpha = \bigcup \alpha < \lambda)(\forall x, y \in T_\alpha)(\text{pr}(x) = \text{pr}(y) \rightarrow x = y)$ .

An  $\omega_1$ -tree is *Aronszajn* if it has no  $\omega_1$ -branches. It is *Souslin* if in addition it has no uncountable antichains.

**THEOREM 1.**

1. (Aronszajn) *There is an Aronszajn tree.*
2. (Gaifman-Specker) *There are  $2^{\aleph_1}$  non-isomorphic Aronszajn trees.*
3. (Jensen) *If  $V = L$ , there is a Souslin tree.*
4. (Jech) *If  $V = L$ , there are  $2^{\aleph_1}$  non-isomorphic Souslin trees.*

For proofs of 1 and 3 we refer the reader to [5]; a proof of 2 may be found in [3], whilst 4 is proved in [4]. It is also proved in [5] that the existence of a Souslin tree is not provable in ZFC alone.

**2. The theorem.** Let  $T$  be an  $\omega_1$ -tree,  $X = \langle X, \leq_X \rangle$  a poset. We say  $T$  is  $X$ -embeddable if there is  $f: T \rightarrow X$  such that  $x <_T y \rightarrow f(x) <_X f(y)$ . We then say  $f$  embeds  $T$  in  $X$ .

**THEOREM 2.** *Let  $T$  be an  $\omega_1$ -tree.*

1.  $T$  is  $X$ -embeddable iff  $T$  is Aronszajn and there are antichains  $A_n$ ,  $n < \omega$ , of  $T$  such that  $T = \bigcup_{n < \omega} A_n$ .
2. If  $T$  is  $R$ -embeddable, then  $T$  is Aronszajn and such that every uncountable subset of  $T$  contains an uncountable antichain of  $T$ .

Proof. 1. If  $f$  embeds  $T$  in  $Q$  then  $T$  is clearly Aronszajn. Also,  $A_q = \{x \in T \mid f(x) = q\}$  is an antichain of  $T$  for each  $q \in Q$ , and  $T = \bigcup_{q \in Q} A_q$ . The converse is trivial.

2. Let  $U \subset T$  be uncountable. Then  $U$  inherits a tree structure from  $T$ . Let  $U^* = \bigcup_{\alpha < \omega_1} U_{\alpha+1}$ . If  $T$  is  $R$ -embeddable, so is  $U$ , whence  $U^*$  is  $Q$ -embeddable. By 1,  $U^*$  is the union of countably many antichains, one of which must be uncountable.

Now, the Aronszajn trees constructed in ZFC are all  $Q$ -embeddable. By the above theorem, no Souslin tree can be  $R$ -embeddable. The question arises, therefore, as to whether there can be Aronszajn trees  $R$ -embeddable but not  $Q$ -embeddable. That such trees cannot be constructed in ZFC follows from the following result, proved in [2]:

**THEOREM 3** (Baumgartner). *If ZFC is consistent, so is ZFC + "Every Aronszajn tree is  $Q$ -embeddable".*

However, the following was announced in [1]:

**THEOREM 4** (Baumgartner). *Assume  $V = L$ . Then there is an Aronszajn tree which is  $R$ -embeddable but not  $Q$ -embeddable.*

Baumgartner's proof, as outlined to us by Richard Laver, involved a forcing construction over initial segments of  $L$ . At the cost of some messy combinatorics, we have adapted this argument to deduce Baumgartner's conclusion from an assumption  $\diamond$ , weaker than  $V = L$ , which is due to R. B. Jensen. This approach allows us to extend Theorem 4 along the lines suggested by Theorem 1.

Let  $\lambda$  be a limit ordinal. A set  $A \subset \lambda$  is *stationary* if it intersects every closed unbounded subset of  $\lambda$ .

**Axiom  $\diamond$ .** There is a sequence  $\langle S_\alpha \mid \alpha < \omega_1 \rangle$  such that  $S_\alpha \subset \alpha$  for each  $\alpha$  and such that whenever  $S \subset \omega_1$ , then  $\{\alpha \in \omega_1 \mid S \cap \alpha = S_\alpha\}$  is stationary in  $\omega_1$ .

**THEOREM 5** (Jensen). *If  $V = L$ , then  $\diamond$  holds.*

Proof. By induction, define  $\langle S_\alpha, C_\alpha \rangle$  as the least pair (under the canonical well-order of  $L$ ) of subsets of  $\alpha$  such that  $C_\alpha$  is closed and unbounded in  $\alpha$  and  $\gamma \in C_\alpha \rightarrow S_\alpha \cap \gamma \neq S_\gamma$ . If no such pair exists, set  $S_\alpha = C_\alpha = \emptyset$ . Assumption that  $\langle S_\alpha \mid \alpha < \omega_1 \rangle$  is not as required now leads speedily to a contradiction. Q.E.D.

We are now ready to prove our theorem. The proof was inspired by arguments in [4].

**THEOREM 6.** *Assume  $\diamond$ . Then there are  $2^{\aleph_1}$  non-isomorphic Aronszajn trees  $R$ -embeddable but not  $Q$ -embeddable.*

Proof. By  $\diamond$  there is a sequence  $\langle h_\alpha \mid \alpha < \omega_1 \rangle$  such that  $h_\alpha: \alpha \rightarrow \alpha$  for each  $\alpha$ , and whenever  $h: \omega_1 \rightarrow \omega_1$ , then  $\{\alpha \in \omega_1 \mid h \upharpoonright \alpha = h_\alpha\}$  is stationary in  $\omega_1$ . We fix this sequence for the rest of this proof.

By induction on  $\alpha < \omega_1$ , for each  $f \in 2^\alpha$  we shall construct an  $(\alpha+1)$ -tree  $T_f$  consisting of sequences of distinct integers of lengths  $\leq \alpha$ . The ordering will be sequence extension. If  $s \in T_f$  and  $\gamma < \text{length}(s)$ , then  $s \upharpoonright \gamma \in T_f$ , whence the height of any  $s$  in  $T_f$  is its length. If  $f, g \in \bigcup_{\alpha < \omega_1} 2^\alpha$  and  $f \subset g$ , then  $T_g$  will be an end-extension of  $T_f$ . Hence for each  $F \in 2^{\omega_1}$ ,  $T(F) = \bigcup_{\alpha < \omega_1} T_{F \upharpoonright \alpha}$  will be an  $\omega_1$ -tree. It will automatically be  $R$ -embeddable.

For, given any  $X \subset \omega$ , let  $f_X \in 2^\omega$  be defined by  $f_X(n) = 1$  iff  $n \in X$ . Then the map  $h: T(F) \rightarrow R$  defined by  $h(s) = f_{\text{range}(s)}$  embeds  $T(F)$  in  $R$ . We shall ensure that no  $T(F)$  is  $Q$ -embeddable, and that  $F, G \in 2^{\omega_1} \ \& \ F \neq G$  implies  $T(F) \not\cong T(G)$ .

As the induction proceeds, we define one-one maps  $\pi_f: T_f \rightarrow \omega_1 - \omega$  so that  $s \subset t \rightarrow \pi_f(s) < \pi_f(t)$ , and so that  $f \subset g \rightarrow \pi_f \subset \pi_g$ .

By  $Q$ , we shall mean  $\omega$  endowed with a dense linear order  $<_Q$ .

We shall carry out the construction so as to preserve the following conditions:

( $\Xi$ ) If  $f \in \bigcup_{\alpha < \omega_1} 2^\alpha$  and  $s \in T_f$ , then  $|\omega - \text{range}(s)| = s_0$ .

( $\Psi$ ) If  $f \in \bigcup_{\alpha < \omega_1} 2^\alpha$  and  $s \in T_f$  and  $x \in [\omega]^{<\omega} - \text{range}(s)$ , there is  $s' \supset s$  on each higher level of  $T_f$  such that  $\text{range}(s') \cap x = \emptyset$ .

Let  $T_\emptyset = \{\emptyset\}$ . Suppose  $\alpha < \omega_1$ ,  $f \in 2^{\alpha+1}$ ,  $T_{f \upharpoonright \alpha}$ ,  $\pi_{f \upharpoonright \alpha}$  are defined, and that  $T_{f \upharpoonright \alpha}$  satisfies ( $\Xi$ ) and ( $\Psi$ ). To obtain  $T_f$ , for each  $s \in \omega^\alpha \cap T_{f \upharpoonright \alpha}$ , add all one-point extensions of  $s$  by distinct integers. This is possible by ( $\Xi$ ),

which guarantees the existence of  $s_0$  such. Clearly,  $T_f$  satisfies  $(\Xi)$  and  $(\Psi)$ . To obtain  $\pi_f$  from  $\pi_{f|a}$ , extend the latter arbitrarily, except for ensuring that if  $g \in 2^{\alpha+1}$  and  $g \neq f$ , then  $\text{range}(\pi_g) \not\cong \text{range}(\pi_f)$ . Since  $2^{\aleph_0} = \aleph_1$  is a trivial consequence of  $\diamond$ , this causes no trouble.

Suppose now that  $\alpha = \bigcup a < \omega_1$ ,  $f \in 2^\alpha$ ,  $T_{f|_\gamma}$ ,  $\pi_{f|_\gamma}$  are defined for all  $\gamma < \alpha$ , and that each  $T_{f|_\gamma}$  satisfies  $(\Xi)$  and  $(\Psi)$ . Let  $T'_f = \bigcup_{\gamma < \alpha} T_{f|_\gamma}$ , an  $\alpha$ -tree. We must decide which  $\alpha$ -branches of  $T'_f$  to extend in order to obtain  $T_f$ . Let  $\pi'_f = \bigcup_{\gamma < \alpha} \pi_{f|_\gamma}$ . Then  $\pi'_f$  induces an  $\alpha$ -tree isomorphic to  $T'_f$  whose elements are countable ordinals. Also,  $s <_{T'_f} t$  implies  $\pi'_f(s) < \pi'_f(t)$ . There are three cases to consider.

Case I.  $h_a$  embeds  $\pi'_f T'_f$  in  $\mathcal{Q}$ . (We understand this to imply that  $\text{domain}(h_a) = \pi'_f T'_f$ .)

Then  $h_a \cdot \pi'_f$  embeds  $T'_f$  in  $\mathcal{Q}$ . (Note that  $f$  is uniquely determined by  $a$  here. For suppose  $g \in 2^\alpha$ ,  $g \neq f$ . Since  $\alpha = \bigcup a$  there is  $\gamma < \alpha$  with  $f|_{\gamma+1} \neq g|_{\gamma+1}$ . By construction,  $\text{range}(\pi_{f|_{\gamma+1}}) \neq \text{range}(\pi_{g|_{\gamma+1}})$ . So, as  $\pi'_f$  and  $\pi'_g$  are order-preserving,  $\text{range}(\pi'_f) \neq \text{range}(\pi'_g)$ . But  $\text{range}(\pi'_f) = \text{domain}(h_a)$ . Thus  $\text{range}(\pi'_g) \neq \text{domain}(h_a)$ , whence  $h_a$  does not embed  $\pi'_g T'_g$  in  $\mathcal{Q}$ .) Let

$$X(a) = \{(s, x) \mid s \in T'_f \ \& \ x \in [\omega]^{<\omega} \ \& \ \text{range}(s) \cap x = \emptyset\}.$$

For  $(s, x), (t, y) \in X(a)$ , say  $(s, x) \leq_a (t, y)$  iff  $s \subset t$  &  $x \subset y$ . This defines a partial order on  $X(a)$ .

Recall that if  $P$  is a poset, a set  $U \subset P$  is *cofinal* if

$$(\forall p \in P)(\exists q \in U)(p \leq_P q).$$

For each  $n \in \omega$ , set

$$A_n^\alpha = \{(s, x) \in X(a) \mid h_a \cdot \pi'_f(s) \geq_Q n \text{ or else}$$

$$(\forall (t, y) \in X(a))[(t, y) \geq_a (s, x) \rightarrow h_a \cdot \pi'_f(t) <_Q n]\}.$$

Clearly, each  $A_n^\alpha$  is cofinal in  $X(a)$ .

Let  $s \in T'_f$ ,  $x \in [\omega]^{<\omega} - \text{range}(s)$ . Let  $\langle a_n \mid n < \omega \rangle$  be cofinal in  $a$  with  $\alpha_0 = \text{length}(s)$ . By  $(\Psi)$  we can find  $s'_0 \in T'_f$ ,  $s'_0 \supset s$ , such that  $\text{length}(s'_0) \geq \alpha_0$  and  $\text{range}(s'_0) \cap x = \emptyset$ . Since  $(s'_0, x) \in X(a)$  and  $A_0^\alpha$  is cofinal, we can find a pair  $(s_0, x_0) \geq_a (s'_0, x)$  in  $A_0^\alpha$ . Let  $m_0 \in \omega - [\text{range}(s_0) \cup x_0]$ , by  $(\Xi)$ . Let  $x'_1 = x_0 \cup \{m_0\}$ . By  $(\Psi)$  we can find  $s'_1 \in T'_f$ ,  $s'_1 \supset s_0$ , such that  $\text{length}(s'_1) \geq \alpha_1$  and  $\text{range}(s'_1) \cap x'_1 = \emptyset$ . Since  $(s'_1, x'_1) \in X(a)$  and  $A_1^\alpha$  is cofinal, we can find  $(s_1, x_1) \geq_a (s'_1, x'_1)$  in  $A_1^\alpha$ . Pick  $m_1 \in \omega - [\text{range}(s_1) \cup x_1]$ , set  $x'_2 = x_1 \cup \{m_1\}$ , and proceed inductively.

Let  $s(x) = \bigcup_{n < \omega} s_n$ . Then  $s(x)$  is an  $\alpha$ -sequence of distinct integers which defines an  $\alpha$ -branch of  $T'_f$ . Also,  $s(x) \supset s$  and  $\text{range}(s(x)) \cap x = \emptyset$ .

Finally, since  $\{m_0, m_1, \dots\} \cap \text{range}(s(x)) = \emptyset$ , we have  $|\omega - \text{range}(s(x))| = \aleph_0$ . We may thus let the  $\alpha$ th level of  $T_f$  consist of one such  $s(x)$  for each pair  $(s, x)$  as above. Then  $T_f$  is an  $(\alpha+1)$ -tree satisfying  $(\Xi)$  and  $(\Psi)$ . Extend  $\pi'_f$  to  $\pi_f$  on  $T_f$  arbitrarily.

Case II. For some  $g \in 2^\alpha$  with  $g \neq f$ ,  $h_a: \pi'_f T'_f \cong \pi'_g T'_g$ . As before,  $f$  and  $g$  are uniquely determined by  $a$ . Let  $s \in T'_f$ ,  $x \in [\omega]^{<\omega} - \text{range}(s)$ . Take  $\langle a_n \mid n < \omega \rangle$  as before. By  $(\Psi)$  we can find  $s_0 \in T'_f$ ,  $s_0 \supset s$ , such that  $\text{length}(s_0) \geq \alpha_0$  and  $\text{range}(s_0) \cap x = \emptyset$ . Let  $m_0 \in \omega - [\text{range}(s_0) \cup x]$ . Let  $x_0 = x \cup \{m_0\}$ . By  $(\Psi)$  again we can find  $s_1 \in T'_f$ ,  $s_1 \supset s_0$ , such that  $\text{length}(s_1) \geq \alpha_1$  and  $\text{range}(s_1) \cap x_0 = \emptyset$ . Let  $m_1 \in \omega - [\text{range}(s_1) \cup x_0]$ , put  $x_1 = x_0 \cup \{m_1\}$ , and proceed inductively. This yields an  $\alpha$ -sequence  $s(x) \supset s$  which determines an  $\alpha$ -branch of  $T'_f$ . Also,  $\text{range}(s(x)) \cap x = \emptyset$  and  $|\omega - \text{range}(s(x))| = \aleph_0$ . Let the  $\alpha$ th level of  $T_f$  consist of one such  $s(x)$  for each such pair  $s, x$ . Similarly for  $T_g$ . The only cause for concern now is if  $\pi'_g^{-1} \cdot h_a \cdot \pi'_f$  extends to an isomorphism of  $T_f$  and  $T_g$ . If it does, pick any distinct  $\alpha$ -branch  $t$  of  $T'_f$  with  $|\omega - \text{range}(t)| = \aleph_0$  and put  $t$  into the  $\alpha$ th level of  $T_f$ . To find such a  $t$ , proceed much as before, but miss the (countably many) branches which already extend. Thus  $\pi'_g^{-1} \cdot h_a \cdot \pi'_f$  cannot now extend to an isomorphism of  $T_f$  and  $T_g$ .

In either case,  $T_f$  and  $T_g$  are  $(\alpha+1)$ -trees satisfying  $(\Xi)$  and  $(\Psi)$ . Extend  $\pi'_f$  and  $\pi'_g$  to  $\pi_f$  and  $\pi_g$  arbitrarily.

Case III. Neither of cases I or II occurs. As in Case II, add one  $s(x)$  for each pair  $s, x$  to obtain  $T_f$ , and extend  $\pi'_f$  arbitrarily.

For  $F \in 2^{\omega_1}$ , set  $T(F) = \bigcup_{\alpha < \omega_1} T_{F|_\alpha}$ , an  $\omega_1$ -tree embeddable in  $\mathbf{R}$ . Let  $\pi_F = \bigcup_{\alpha < \omega_1} \pi_{F|_\alpha}$ . Then  $\pi_F: T(F) \rightarrow \omega_1 - \omega$  and  $s <_{T(F)} t \rightarrow \pi_F(s) < \pi_F(t)$ . Also,  $\pi_F$  induces a tree isomorphic to  $T(F)$  whose elements are countable ordinals.

Suppose  $T(F)$  is  $\mathcal{Q}$ -embeddable. Then there is an  $h$  which embeds  $\pi_F T(F)$  in  $\mathcal{Q}$ . Let

$$\begin{aligned} A = \{ & a \in \omega_1 \mid a = \bigcup \alpha \cdot \& \cdot [\pi_F T(F)] \upharpoonright \alpha = \pi'_F T'_{F|_\alpha} \cdot \& \cdot \\ & \cdot \& \cdot h \upharpoonright \alpha \text{ embeds } \pi'_F T'_{F|_\alpha} \text{ in } \mathcal{Q} \cdot \& \cdot \\ & \cdot \& \cdot (\forall s \in T'_{F|_\alpha})(\forall x \in [\omega]^{<\omega} - \text{range}(s)) (\forall n >_Q h(\pi_F(s))) \\ & [(\exists t \in T(F))(t \supset s \ \& \ \text{range}(t) \cap x = \emptyset \ \& \ h(\pi_F(t)) \geq_Q n) \\ & \rightarrow (\exists t \in T'_{F|_\alpha})(t \supset s \ \& \ \text{range}(t) \cap x = \emptyset \ \& \ h(\pi_F(t)) \geq_Q n)] \}. \end{aligned}$$

It is easily seen that  $A$  is closed and unbounded in  $\omega_1$ . Hence by  $\diamond$  there is  $a \in A$  such that  $h \upharpoonright a = h_a$ . Thus, by the definition of  $A$ , Case I applied in constructing  $T_{F|_\alpha}$  from  $T'_{F|_\alpha}$ . Let  $s \in \omega^\alpha \cap T(F)$ . Let  $n = h(\pi_F(s))$ . By construction, let  $(t, y) \in A_n^\alpha$  be such that  $\text{range}(s) \cap y = \emptyset$  and  $t \subset s$ .



As  $h$  is order-preserving,  $h(\pi_F(t)) <_Q n$ . By the definition of  $A$  and the relation of  $s$  to  $t$ , there is  $t' \in T'_{F|a}$  such that  $t' \supset t$  and  $\text{range}(t') \cap y = \emptyset$ , and such that  $h(\pi_F(t')) \geq_Q n$ . In particular,  $(t, y), (t', y) \in X(a)$  and  $(t', y) \geq_a (t, y)$ . But look,  $(t, y) \in A_n^a$ , so by definition we must have that  $h_a(\pi_F(t)) \geq_Q n$ . Since  $h \upharpoonright a = h_a$ ; this contradicts our earlier inequality. Hence  $T(F)$  cannot be  $Q$ -embeddable.

Suppose now that for some  $F, G \in 2^{\omega_1}$ ,  $F \neq G$ , we have  $T(F) \cong T(G)$ . Let  $h: \pi_F'' T(F) \cong \pi_G'' T(G)$ . Pick  $\alpha_0 < \omega_1$  such that  $F \upharpoonright \alpha_0 \neq G \upharpoonright \alpha_0$ . Let

$$A = \{a \in \omega_1 \mid a = \bigcup \alpha > \alpha_0 \cdot \& \cdot [\pi_F'' T(F)] \upharpoonright \alpha = \pi'_{F|a}'' T'_{F|a} \cdot \& \cdot [\pi_G'' T(G)] \upharpoonright \alpha = \pi'_{G|a}'' T'_{G|a} \cdot \& \cdot h \upharpoonright \alpha: \pi'_{F|a}'' T'_{F|a} \cong \pi'_{G|a}'' T'_{G|a}\}.$$

Clearly,  $A$  is closed and unbounded in  $\omega_1$ . By  $\diamond$ , there is  $a \in A$  such that  $h \upharpoonright a = h_a$ . Thus Case II applied in constructing  $T'_{F|a}$  from  $T'_{F|a}$  and  $T'_{G|a}$  from  $T'_{G|a}$ . This means that the map  $\pi'_{G|a} \cdot h_a \cdot \pi'^{-1}_{F|a}$  does not extend to an isomorphism of  $T'_{F|a}$  and  $T'_{G|a}$ , which is absurd, since  $\pi'^{-1}_{G|a} \cdot h \cdot \pi_F$  extends it. Thus  $T(F)$  and  $T(G)$  are not isomorphic. The proof is complete.

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## Shapes of finite-dimensional compacta

by

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**1. Introduction.** The results of this paper deal with shapes of finite-dimensional compact metric spaces (see [4] for definitions concerning the concept of shape). In Theorem 1 below we give a characterization of shapes of finite-dimensional compact metric spaces (i.e. *compacta*) in terms of embeddings in Euclidean  $n$ -space  $E^n$ . In an earlier paper the author obtained a characterization of shapes of compacta (with no dimensional restriction) in terms of embeddings in the Hilbert cube [8]. In a sense the results obtained here are motivated by [8], and to some extent the general structure of the proof of Theorem 1 is a modification of the argument used in [8], but the present paper does not involve any infinite-dimensional topology. For the sake of completeness we give a short summary of the infinite-dimensional characterization at the end of the Introduction. We use the notation  $\text{Sh}(X) = \text{Sh}(Y)$  to indicate that compacta  $X$  and  $Y$  have the same shape.

**THEOREM 1.** *Let  $X, Y$  be compacta such that  $\dim X, \dim Y \leq m$ .*

(a) *For any integer  $n \geq 2m + 2$  there exist copies  $X', Y' \subset E^n$  (of  $X, Y$  respectively) such that if  $\text{Sh}(X) = \text{Sh}(Y)$ , then  $E^n \setminus X'$  and  $E^n \setminus Y'$  are homeomorphic.*

(b) *For any integer  $n \geq 3m + 3$  there exist copies  $X', Y' \subset E^n$  (of  $X, Y$  respectively) such that if  $E^n \setminus X'$  and  $E^n \setminus Y'$  are homeomorphic, then  $\text{Sh}(X) = \text{Sh}(Y)$ .*

We remark that a similar result holds for embeddings of  $X$  and  $Y$  in the  $n$ -sphere  $S^n$ .

For prerequisites we will need some elementary facts concerning the piecewise-linear topology of  $E^n$  plus an isotopy extension theorem from [11]. We also use a characterization of dimension in terms of mappings onto polyhedra in  $E^n$  (see [14], p. 111). As for techniques we remark that part (a) of Theorem 1 is the most difficult part of the proof. Roughly the idea is to construct a sequence  $\{h_i\}_{i=1}^{\infty}$  of homeomorphisms

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