

## Embedding of Boolean algebras in $P(\omega)/\text{fin}$

by

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**Abstract.** A model of ZFC set theory is constructed in which  $c$  (the continuum) is large, MA ( $\sigma$ -linked) holds and each compact 0-dimensional space of weight  $c$  is a continuous image of  $\omega^*$ .

A well-known theorem of Parovichenko (see e.g. [C-N]) says that, if CH (the continuum hypothesis) is assumed, then each Boolean algebra of cardinality  $\leq c = 2^\omega$  can be embedded into  $P(\omega)/\text{fin}$ . What happens if CH fails? Kunen [K] proved that after adding  $\omega_2$  Cohen reals to a model of CH, there is an algebra of cardinality  $c$  which cannot be embedded into  $P(\omega)/\text{fin}$ . In [F] a model of ZFC + MA is constructed in which  $c = 2^{\omega_1}$  and the powerset algebra  $P(\omega_1)$  cannot be embedded into  $P(\omega)/\text{fin}$ . There is also a model in which  $c = \omega_2$  and the measure algebra does not embed into  $P(\omega)/\text{fin}$ . On the other hand, Laver shows (in [L]) that, consistently with a large value for  $c$ , each linear ordering of cardinality  $\leq c$  embeds in  $\omega^\omega$  (ordered by eventual dominance)<sup>(1)</sup>.

Using Laver's technique we prove here the following:

**THEOREM.** *There is a model of ZFC set theory with arbitrarily large  $c$  in which each Boolean algebra of cardinality  $\leq c$  can be embedded into  $P(\omega)/\text{fin}$ . In addition, the Martin Axiom for  $\sigma$ -linked orderings holds in the model.*

In particular, it follows that the conclusion of the Parovichenko theorem is not equivalent to CH.<sup>1</sup>

**Proof.** In the ground model  $V$  we fix a regular cardinal  $\kappa > \omega_1$ , with  $2^{<\kappa} = \kappa$ , and force with a  $\kappa$ -stage finite support iteration  $\langle P_\alpha, Q_\alpha; \alpha < \kappa \rangle$  so that the resulting extension  $V[G]$  will have the required properties.

In a usual way choose names  $Q_\alpha$ ,  $\alpha < \kappa$  and  $\alpha$  odd, so that

$$P_\alpha \Vdash "Q_\alpha \text{ is } \sigma\text{-linked}"$$

and  $V[G] \models \text{MA} (\sigma\text{-linked})$ .

<sup>(1)</sup> The consistency of ZFC +  $\neg\text{CH}$  + each Boolean algebra of size  $\leq c$  can be embedded into  $P(\omega)/\text{fin}$  has been pointed out without proof by Baumgartner in his article *Application of the proper forcing axiom in Handbook of Set-Theoretic Topology*.

At even stages  $\alpha$ ,  $Q_\alpha$  adds a subset  $A_\alpha \subseteq \omega$  which fills a certain gap in  $P(\omega)/\text{fin}$ . To be more precise, let  $C(\alpha)$  consist of finite zero-one functions  $s$  with  $\text{dm}(s) \subseteq \alpha$ . Set

$$D(\alpha) = \{s \in C(\alpha) : \forall \xi \in \text{dm}(s) \ \xi \text{ is even}\}.$$

Each  $s \in D(\alpha)$  determines a Boolean combination

$$A(s) = \bigcap_{s(\xi)=0} A_\xi \cap \bigcap_{s(\eta)=1} (\omega \setminus A_\eta).$$

Now, in  $V^{P_\alpha}$  a pair  $\langle S, T \rangle$  is called a *generic gap* if  $S, T \subseteq D(\alpha)$  and the following two properties hold:

(1)  $A(s) \cap A(t) \in \text{fin}$ , for each  $s \in S$  and  $t \in T$ .

(2) If  $A(s) \subseteq_* A(s_1) \cup \dots \cup A(s_n)$ , for some  $s \in D(\alpha)$  and  $s_1, \dots, s_n \in S$ , then  $s \in S$ ; similarly for  $T$ .

Standard reasoning shows that there is a sequence  $\langle S_\alpha, T_\alpha \rangle$  such that  $\langle S_\alpha, T_\alpha \rangle$  is a  $P_\alpha$ -name of a generic gap in  $V^{P_\alpha}$  and each  $\kappa$ -generated gap occurs in the sequence, for arbitrarily large  $\alpha < \kappa$ .

Now,  $Q_\alpha$  consists of  $q = \langle x_q, a_q, y_q, m_q \rangle$ , where  $x_q \subseteq S_\alpha$ ,  $y_q \subseteq T_\alpha$  are finite,  $a_q$  is a zero-one sequence of length  $m_q$  and

$$\bigcup_{s \in x_q} A(s) \cap \bigcup_{t \in y_q} A(t) \subseteq m_q.$$

The ordering on  $Q_\alpha$  is defined thus:

$$p \leq q \quad \text{iff} \quad x_p \supseteq x_q, a_p \supseteq a_q, y_p \supseteq y_q, m_p \geq m_q$$

and for each  $i$  with  $m_q \leq i < m_p$  we have

$$a_p(i) = 0, \quad \text{for } i \in \bigcup \{A(s) : s \in x_q\},$$

$$a_p(i) = 1, \quad \text{for } i \in \bigcup \{A(t) : t \in y_q\}.$$

If  $H \subseteq Q_\alpha$  is a generic filter, then

$$A_\alpha = \{i \in \omega : \exists q \in H [a_q(i) = 0]\}$$

fills the gap  $\langle S_\alpha, T_\alpha \rangle$ , i.e. we have

$$A_s \subseteq_* A_\alpha, \quad \text{for each } s \in S_\alpha,$$

$$A_t \cap A_\alpha \in \text{fin}, \quad \text{for each } t \in T_\alpha.$$

**Remark.** Let us note the following property of  $Q_\alpha$ . Let  $s \in D(\alpha)$ . If  $A(s) \subseteq_* A_\alpha$ , then  $s \in S_\alpha$ . Similarly, if  $A(t) \cap A_\alpha =_* \emptyset$ , then  $t \in T_\alpha$ . Indeed,  $q \Vdash A(s) \cap n \subseteq A_\alpha$  implies

$$A(s) \setminus \max\{n, m_q\} \subseteq \bigcup \{A(s') : s' \in x_q\}$$

for otherwise we have a  $q' \leq q$  and an  $i \in A(s) \setminus n$  with  $a_{q'}(i) = 1$ . Thus,  $s \in S_\alpha$ , by the condition (2) above.

It remains to show that  $P_\alpha$  has the c.c.c. Indeed, assuming that for the moment, it is clear that  $c = \kappa$  and MA ( $\sigma$ -linked) hold in the resulting model  $V[G]$ . Suppose that

$B = \{b_\alpha : \alpha < \kappa\}$  is a Boolean algebra and let  $b_\alpha$  be such that  $B_{\alpha+1}$  is generated by  $B_\alpha$  and  $b_\alpha$ , for  $\alpha < \kappa$ . In view of the Sikorski extension theorem it suffices to find a function  $f : \kappa \rightarrow \kappa$  (with even values) such that for each  $s \in C(\kappa)$  we have

$$(*) \quad b(s) = 0 \quad \text{iff} \quad A(s') \in \text{fin},$$

(where  $b(s) = (\prod_{s(\xi)=0} b_\xi) \cdot (\prod_{s(\eta)=1} b_\eta)$ ) and  $s^f(f(\beta)) = s(\beta)$ , for  $\beta \in \text{dm}(s)$ ).

Assume that  $f$  has already been defined for all  $\beta < \alpha$  so that  $(*)$  holds for each  $s \in C(\alpha)$  and  $b_\alpha$  not in the subalgebra generated by  $\{b_\beta : \beta < \alpha\}$ . Let

$$S = \{s' : s' \in C(\alpha) \text{ and } b(s) \leq b_\alpha\},$$

$$T = \{s' : s' \in C(\alpha) \text{ and } b(s) \cdot b_\alpha = 0\}.$$

Now,  $\langle S, T \rangle$  generates a generic gap  $\langle S_\gamma, T_\gamma \rangle$  at some stage  $\gamma > \sup \{f(\beta) : \beta < \alpha\}$  and we define  $f(\alpha) = \gamma$ . To check  $(*)$ , let  $\max \text{dm}(s) = \alpha$  and for example  $s(\alpha) = 1$ . Then  $b(s) = 0$  implies  $b(s|\alpha) \leq b_\alpha$  and hence  $(s|\alpha)' \in S$ , which gives  $A((s|\alpha)') \subseteq_* A_\gamma$ , i.e.  $A(s') =_* \emptyset$ .

Conversely, from  $A(s') \in \text{fin}$  follows  $A((s|\alpha)') \subseteq_* A_\gamma$  and hence  $(s|\alpha)' \in S$ , by the Remark above. This means

$$A((s|\alpha)') \subseteq_* A(s'_1) \cup \dots \cup A(s'_l)$$

for some  $s'_1, \dots, s'_l \in S$ . By the inductive assumption

$$b(s|\alpha) \leq b(s_1) + \dots + b(s_l).$$

Since  $b(s_1), \dots, b(s_l) \leq b_\alpha$  we obtain  $b(s|\alpha) \leq b_\alpha$  and hence  $b(s) = 0$ .

The rest of the paper contains the proof of the countable chain condition.

Let  $E_\alpha \subseteq P_\alpha$  consist of all  $p \in P_\alpha$  having the following properties:

(3) For each even  $\xi \in \text{supp}(p)$ , there are  $x_\xi, a_\xi, y_\xi$  such that  $p \Vdash \xi \Vdash p(\xi) = \langle x_\xi, x_\xi, y_\xi \rangle$ , for each  $s \in x_\xi \cup y_\xi$  with  $\text{dm}(s) \subseteq \text{supp}(p)$ .

(4) There is a number  $l = l(p)$  such that for each even  $\xi \in \text{supp}(p)$ , we have  $l(a_\xi) = l(p)$  (i.e. all the  $a$ 's are of the same length).

For odd  $\alpha$  we have  $P_\alpha \Vdash "Q_\alpha \text{ is } \sigma\text{-linked}"$ , and hence we can choose a  $P_\alpha$ -name  $h_\alpha$  such that

$$P_\alpha \Vdash "h_\alpha : Q_\alpha \rightarrow \omega \text{ and } \forall n \ h^{-1}(n) \text{ is linked}."$$

Let  $E_\alpha^*$  consist of all  $p \in E_\alpha$  which in addition satisfy

(5) For each odd  $\xi \in \text{supp}(p)$ , there is an  $n < \omega$  such that

$$p \Vdash \xi \Vdash "h_\xi(p(\xi)) = n".$$

**LEMMA 1.** For each  $p \in P_\alpha$  and  $m \in \omega$ , there is a  $q \leq p$  such that  $q \in E_\alpha^*$  and  $l(q) \geq m$ .

**Proof.** By induction on  $\alpha$ . Since supports are finite, it suffices to consider only the case  $\alpha = \beta + 1$  and  $\beta \in \text{supp}(p)$ . If  $\beta$  is odd, we find a  $q' \leq p \upharpoonright \beta$  and an  $n \in \omega$  with  $q' \Vdash "h_\beta(p(\beta)) = n"$ . By the inductive assumption, there is a  $q'' \leq q'$  such that  $q'' \in E_\beta^*$  and

$l(q'') \geq m$ . Now, if  $q \in P_\alpha$  is defined by  $q|\beta = q''$  and  $q(\beta) = p(\beta)$ , then  $q \leq p$ ,  $q \in E_\alpha^*$  and  $l(q) \geq m$ . If  $\beta$  is even, then we find a  $q' \leq p|\beta$  and  $\langle x_\beta, a_\beta, y_\beta \rangle$  so that

$$q' \Vdash "p(\beta) = \langle x_\beta, a_\beta, y_\beta \rangle"$$

and  $dm(s) \subseteq \text{supp}(q')$ , for each  $s \in x_\beta \cup y_\beta$ . By the inductive assumption, there is a  $q'' \in E_\beta^*$  with  $l(q'') \geq \max\{m, l(a_\beta)\}$ . In particular,

$$q'' \Vdash "A(s) \cap A(t) \subseteq l(a_\beta)", \quad \text{for } s \in x_\beta \text{ and } t \in y_\beta,$$

and  $q''$  determines the segments  $A_\xi \cap l(q'')$ , for each  $\xi \in dm(s)$  and  $s \in x_\beta \cup y_\beta$ . Hence, we can define an  $a \supseteq a_\beta$  with  $l(a) = l(q'')$  such that  $q'' \Vdash "\langle x_\beta, a, y_\beta \rangle \leq \langle x_\beta, a_\beta, y_\beta \rangle"$ .

Now, if  $q \in P_\alpha$  is such that  $q|\beta = q''$  and  $q(\beta) = \langle x_\beta, a, y_\beta \rangle$ , then  $q$  has the required properties.

**LEMMA 2.** Let  $p \in E_\alpha$ ,  $s \in D(\alpha)$  and  $\beta = \max dm(s)$ . If  $p \Vdash "A(s) \in \text{fin}"$ , then there is a  $q \leq p$  with  $q \in E_\alpha$  and  $l(q) = l(p)$ , such that if  $q|\beta \Vdash "q(\beta) = \langle x_\beta, a_\beta, y_\beta \rangle"$ , then

$$s|\beta \in x_\beta \quad \text{if } s(\beta) = 1,$$

$$s|\beta \in y_\beta \quad \text{if } s(\beta) = 0.$$

*Proof.* Induction on  $\alpha$ . From the assumptions of the lemma it follows that  $p|(\beta+1) \Vdash "A(s) \in \text{fin}"$  and hence we may assume  $\alpha = \beta+1$ . Suppose

$$p|\beta \Vdash p(\beta) = \langle x, a, y \rangle$$

and let  $s(\beta) = 1$ . Thus, we have

$$p|(\beta+1) \Vdash "A(s|\beta) \subseteq {}_* A_\beta"$$

and hence  $p|\beta \Vdash "s|\beta \in S_\beta"$ . In particular,

$$p|\beta \Vdash "A(s|\beta) \cap A(t) \in \text{fin}", \quad \text{for each } t \in y.$$

We have to show

$$p|\beta \Vdash "A(s|\beta) \cap A(t) \subseteq l(p)", \quad \text{for } t \in y.$$

If  $s|\beta \cup t$  is not a function, then we have

$$P_\beta \Vdash "A(s|\beta) \cap A(t) = \emptyset".$$

If  $s|\beta \cup t$  is a function, then  $A(s|\beta) \cap A(t) = A(s|\beta \cup t)$  and we can apply the inductive assumption: there is a  $q_t \in E_\beta$  with  $q_t \leq p|\beta$  and  $l(q_t) = l(p)$  such that if  $\beta_t = \max dm(s|\beta \cup t)$  and

$$q_t|\beta_t \Vdash "q_t(\beta) = \langle x_t, a_t, y_t \rangle"$$

then  $(s|\beta \cup t)|\beta_t$  is in  $x_t$  or  $y_t$  according to the value at  $\beta_t$ . In either case we obtain

$$q_t \Vdash "A(s|\beta \cup t) \subseteq l(q_t) = l(p)".$$

Repeating this for each  $t \in y$  such that  $s|\beta \cup t$  is a function, we find a  $q' \in E_\beta$  with  $q' \leq p|\beta$ ,  $l(q') = l(p)$  and

$$q' \Vdash "A(s|\beta) \cup A(t) \subseteq l(p)", \quad \text{for each } t \in y.$$

Hence, if  $q|\beta = q'$  and  $q(\beta) = \langle x \cup \{s|\beta\}, a, y \rangle$ , then  $q$  has the required properties. The case  $s(\beta) = 0$  is similar.

**LEMMA 3.** Assume that  $p, q \in E_\alpha^*$ ,  $l(p) = l(q)$  and for each  $\xi \in \text{supp}(p) \cap \text{supp}(q)$ :

(6) if  $\xi$  is odd and  $p|\xi \Vdash "h_\xi(p(\xi)) = n"$ ,  $q|\xi \Vdash "h_\xi(q(\xi)) = m"$ , then  $n = m$ .

(7) if  $\xi$  is even and  $p|\xi \Vdash "p(\xi) = \langle x_\xi, a_\xi, y_\xi \rangle"$ ,  $q|\xi \Vdash "q(\xi) = \langle u_\xi, b_\xi, v_\xi \rangle"$ , then  $a_\xi = b_\xi$ .

Then  $p$  and  $q$  are compatible. In fact, there is an  $r \leq p, q$  with  $r \in E_\alpha$  and  $l(r) = l(p) = l(q)$ .

*Proof.* Induction on  $\alpha$ . The only essential case is  $\alpha = \beta+1$  and  $\beta \in \text{supp}(p) \cap \text{supp}(q)$ .

Suppose that  $\beta$  is odd. By the inductive assumption there is an  $r' \leq p|\beta, q|\beta$  with  $r' \in E_\beta$  and  $l(r') = l(p)$ . Since

$$r' \Vdash "h_\beta(p(\beta)) = h_\beta(q(\beta))"$$

it follows that there is an  $r^*$  such that

$$r' \Vdash "r^* \leq p(\beta), q(\beta)".$$

Now, if  $r|\beta = r'$  and  $r(\beta) = r^*$ , then  $r \in E_\alpha$ ,  $r \leq p, q$  and  $l(r) = l(p)$  (it may happen that  $r \notin E_\alpha^*$ , since  $r'$  need not determine the value  $h_\beta(r^*)$ ).

Let now  $\beta$  be even. Then we have

$$p|\beta \Vdash "p(\beta) = \langle x, a, y \rangle", \quad q|\beta \Vdash "q(\beta) = \langle u, a, v \rangle".$$

By the inductive assumption there is an  $r' \in E_\beta$  with  $r' \leq p|\beta, q|\beta$  and  $l(r') = l(p|\beta)$ . It follows that

$$r' \Vdash "A(s) \cap A(t) \in \text{fin}", \quad \text{for each } s \in x \cup u \text{ and } t \in y \cup v.$$

Now, we may apply Lemma 2 repeatedly to  $r'$  and to each  $A(s \cup t)$  such that  $s \cup t$  is a function. The properties of the resulting  $r^*$  ensure that there is an  $r \in P_\alpha$  such that  $r|\beta = r^*$ ,  $r(\beta) = \langle x \cup u, a, y \cup v \rangle$  and  $r \in E_\alpha$ . In particular,  $r \leq p, q$  and  $l(r) = l(p) = l(q)$ , which finishes the proof of the lemma.

Suppose now that  $C \subseteq P_\alpha$  is an uncountable antichain. By Lemma 1 we may assume  $C \subseteq E_\alpha^*$ . Applying the  $\Delta$ -system lemma to  $\{\text{supp}(p) : p \in C\}$  and some elementary operations we find an uncountable  $C_0 \subseteq C$  so that any  $p, q \in C_0$  are as in Lemma 3, a contradiction. Thus, each  $P_\alpha$  has the c.c.c. and the proof is complete.

**Remark.** It is not difficult to see that, in the model constructed, each automorphism of a given algebra  $B$  can be, after embedding, extended to an automorphism of  $P(\omega)/\text{fin}$ .

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## On bounded paradoxical subsets of the plane

by

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**Abstract.** We give a precise lower bound for the number of pieces required in a bounded paradoxical subset of the plane.

Intuitively, a subset  $X$  of a metric space is said to be *paradoxical* if it admits a partition  $X = A \cup B$  such that each of the sets  $A, B$  can be subdivided into finitely many pieces which can be reassembled via isometries to produce  $X$ ; if  $A$  is subdivided into  $m$  pieces and  $B$  is subdivided into  $n$  pieces, the set  $X$  is said to be  $(m, n)$ -paradoxical.

The Sierpiński-Mazurkiewicz paradox is that there is a  $(1,1)$ -paradoxical subset of the plane [MS]. Hadwiger, Debrunner and Klee [HDK, p. 80] have shown that a *bounded*  $(m, n)$ -paradoxical subset of the plane must satisfy  $m+n > 2$ . A bounded  $(1, 3)$ -paradoxical subset of the plane has recently been constructed by Just [J].

Our main purpose here is to show that there is no bounded  $(1,2)$ -paradoxical subset of the plane. This improves the result of Hadwiger, Debrunner and Klee, and renders optimal the recent construction of Just. We also construct here a bounded  $(2,2)$ -paradoxical subset of the plane.

**DEFINITION 1.**  $X$  is an  $(m, n)$ -paradoxical subset of the plane if  $X$  is nonempty, and there are subsets  $C_1, \dots, C_m, D_1, \dots, D_n$  of  $X$  and planar isometries  $G_1, \dots, G_m, H_1, \dots, H_n$ , such that  $P_1 = \{C_i\}$ ,  $P_2 = \{D_j\}$  and  $P_3 = \{G_i(C_i)\} \cup \{H_j(D_j)\}$  are each partitions of  $X$ .

**DEFINITION 2.** Let  $X$  be an  $(m, n)$ -paradoxical subset of the plane whose paradoxical decomposition is witnessed by subsets  $C_1, \dots, C_m, D_1, \dots, D_n$  and planar isometries  $G_1, \dots, G_m, H_1, \dots, H_n$ . Write  $\mathcal{C} = \{C_1, \dots, C_m\}$ ,  $\mathcal{D} = \{D_1, \dots, D_n\}$ ,  $\mathcal{G} = \{G_1, \dots, G_m\}$ , and  $\mathcal{H} = \{H_1, \dots, H_n\}$ . We define the *associated directed graph*  $\Gamma = \Gamma(\mathcal{C}, \mathcal{D}, \mathcal{G}, \mathcal{H})$  of the decomposition.  $\Gamma$  is an infinite directed graph with vertex set  $V(\Gamma) = X$ . The set of darts (i.e. directed edges) of  $\Gamma$  consists of all pairs  $(x, G_i(x))$  and  $(x, H_j(x))$ , where  $x \in C_i \cap D_j$ .

It is helpful, when drawing diagrams, to label each dart of  $\Gamma$  with the planar isometry that determined its second coordinate.

Observe that every  $x$  in  $V(\Gamma)$  has invalency 1 and outvalency 2.

**LEMMA 1.** *Let  $\Gamma$  be an infinite directed graph with invalency 1 at each vertex, and suppose furthermore that  $\Gamma$  is connected. Then  $\Gamma$  contains at most one cycle.*