

Elementary methods in the theory of L -functions, VII Upper bound for $L(1, \chi)$

by

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1. If χ is a real nonprincipal character (mod D), then the upper bound which one can give for $L(1, \chi)$ is closely connected with the upper bound of

$$(1.1) \quad S_x = \max_{1 \leq a < b \leq D} \left| \sum_{n=a}^b \chi(n) \right|.$$

Using the trivial $S_x \leq D$ one can easily prove $L(1, \chi) \leq \log D + O(1)$; by means of the Pólya–Vinogradov inequality $S_x \leq c \sqrt{D} \log D$

$$(1.2) \quad L(1, \chi) \leq \left(\frac{1}{2} + o(1)\right) \log D$$

(see Pólya [7]) can be proved.

If $D = p$ is a prime, χ a real nonprincipal character (mod p), then making use of Burgess's inequality [1]

$$(1.3) \quad \left| \sum_{n=N+1}^{N+H} \chi(n) \right| \leq \varepsilon H \quad \text{for } H > p^{1/4+\varepsilon}, p > p_0(\varepsilon),$$

S. Chowla [4] in 1964 proved the inequality

$$(1.4) \quad L(1, \chi_p) \leq \left(\frac{1}{2} + o(1)\right) \log p.$$

Burgess [2] showed in 1966 that

$$(1.5) \quad L(1, \chi_p) < 0.2456 \dots \log p.$$

Wirsing (unpublished) improved it to

$$(1.6) \quad L(1, \chi_p) < \frac{1}{2} (\sqrt{2} - 1 + o(1)) \log p \approx 0.207 \log p.$$

P. J. Stephens [8] showed in 1972 using a method of Wirsing that

$$(1.7) \quad L(1, \chi_p) < \frac{1}{2} \left(1 - \frac{1}{\sqrt{e}} + o(1)\right) \log p \approx 0.197 \log p.$$

Now we give an elementary proof of Stephens's result (using Burgess's inequality) generalizing it for real primitive characters, whose modulus is not necessarily prime, and improve (1.2) for real non-principal characters. Our result will follow from the following general theorem.

THEOREM 1. *If θ is a completely multiplicative function, which takes only the values $+1, 0, -1$, x a real number for which*

$$(1.8) \quad \sum_{n \leq x} \theta(n) \leq \varepsilon x$$

then

$$(1.9) \quad \sum_{d \leq x} \frac{\theta(d)}{d} \leq 2 \left(1 - \frac{1}{\sqrt{e}} + \delta \right) \log x$$

where $\delta = \delta(\varepsilon, x) \rightarrow 0$ if $x \rightarrow \infty$ and $\varepsilon \rightarrow 0$.

Theorem 1 is the best possible, because if we choose

$$\theta(p) = \begin{cases} 1 & \text{for } p \leq x^{1/\sqrt{e}}, \\ -1 & \text{for } p > x^{1/\sqrt{e}} \end{cases} \quad (p \text{ is a prime})$$

then it is easy to see that (1.8) is true with $\varepsilon = o(1)$ and that in (1.9) equality holds with $\delta = o(1)$.

But Burgess [3] proved that if χ is a nonprincipal character (mod D), then

$$(1.10) \quad \left| \sum_{n=N+1}^{N+H} \chi(n) \right| \leq \varepsilon H \quad \text{for } H \geq D^{\tau_x + \varepsilon}, D > D_0(\varepsilon)$$

where if χ is a primitive character then $\tau_x = 1/4$ and for an arbitrary χ , $\tau_x = 3/8$. Thus using (1.10) we have by partial summation

$$(1.11) \quad \left| \sum_{D^{\tau_x + \varepsilon} < d \leq D} \frac{\chi(d)}{d} \right| \leq \varepsilon \log D$$

and using the trivial estimation $S_x \leq D$ by means of Abel's inequality we get

$$(1.12) \quad \left| \sum_{d > D} \frac{\chi(d)}{d} \right| \leq 1.$$

So using Theorem 1 with $x = D^{\tau_x + \varepsilon}$ we have from (1.11) and (1.12)

THEOREM 2. *If χ is a real primitive character (mod D), then*

$$(1.13) \quad L(1, \chi) \leq \frac{1}{2} \left(1 - \frac{1}{\sqrt{e}} + o(1) \right) \log D$$

if χ is a real nonprincipal character (mod D), then

$$(1.14) \quad L(1, \chi) \leq \frac{3}{4} \left(1 - \frac{1}{\sqrt{e}} + o(1) \right) \log D.$$

(1.13) is in the following sense the best possible for $D = p$: If the least quadratic non-residue (mod p),

$$N(p) \geq p^{\frac{1}{4\sqrt{e}} - o(1)}$$

then it is easy to see that in (1.13) the equality is valid. Thus any improvement of (1.13) is only possible if we improve Burgess's theorem [1] concerning the least quadratic non-residue (mod p) to

$$(1.15) \quad N(p) = O(p^\eta)$$

with an $\eta < \frac{1}{4\sqrt{e}}$.

The upper bound of $L(1, \chi)$ is in connection with the class number and fundamental unit of quadratic fields. Using (1.4) S. Chowla ([4], [5]) proved that if p is a prime $\equiv 1 \pmod{4}$, then for the class number $h(p)$, and fundamental unit $\varepsilon > 1$ of $Q(\sqrt{p})$ one has

$$(1.16) \quad h(p) \leq \left(\frac{1}{4} + o(1) \right) \log p$$

and

$$(1.17) \quad \varepsilon \leq e^{\left(\frac{1}{8} + o(1) \right) \sqrt{p} \log p}.$$

He also proved [5] that if p is a prime $\equiv 3 \pmod{4}$, then for the class number $h(-p)$ of $Q(\sqrt{-p})$

$$(1.18) \quad h(-p) \leq \left(\frac{1}{4\pi} + o(1) \right) \sqrt{p} \log p$$

holds. If D or $-D$, respectively is not a prime but a fundamental discriminant the best known upper bounds for class numbers of quadratic fields belonging to the discriminant D or $-D$ respectively, are due to Landau [6], who proved the inequalities

$$(1.19) \quad h(D) \leq \left(\frac{1}{2} + o(1) \right) \sqrt{D} \quad (D > 0)$$

and

$$(1.20) \quad h(-D) \leq \left(\frac{1}{2\pi} + o(1) \right) \sqrt{D} \log D \quad (-D < 0).$$

Taking into account the well-known class number formulae

$$(1.21) \quad 2h(D)\log \varepsilon = \sqrt{D}L(1, \chi) \quad (\chi(n) = \left(\frac{D}{n}\right), D > 0),$$

$$(1.22) \quad h(-D) = \frac{\sqrt{D}}{\pi}L(1, \chi) \quad (\chi(n) = \left(\frac{-D}{n}\right), -D < -4)$$

and the inequality

$$(1.23) \quad \varepsilon \geq \frac{1}{2}(\sqrt{D} + 1),$$

Theorem 2 gives the following improvements of the results of S. Chowla and Landau ((1.16)–(1.20)):

THEOREM 3. For the class number $h(D)$ and for the fundamental real unit $\varepsilon > 1$ of the real quadratic field belonging to the fundamental discriminant $D > 0$ the inequalities

$$(1.24) \quad h(D) \leq \frac{1}{2} \left(1 - \frac{1}{\sqrt{e}} + o(1)\right) \sqrt{D}$$

and

$$(1.25) \quad \varepsilon \leq e^{\frac{1}{4} \left(1 - \frac{1}{\sqrt{e}} + o(1)\right) \sqrt{D} \log D}$$

hold.

THEOREM 4. For the class number $h(-D)$ of the imaginary quadratic field belonging to the fundamental discriminant $-D < 0$ the inequality

$$(1.26) \quad h(-D) \leq \frac{1}{2\pi} \left(1 - \frac{1}{\sqrt{e}} + o(1)\right) \sqrt{D} \log D$$

holds.

2. To prove Theorem 1 first we note that if

$$(2.1) \quad g(n) = \sum_{d|n} \theta(d)$$

then as $\theta(d) = O(1)$, we get

$$(2.2) \quad \sum_{n \leq x} g(n) = \sum_{d \leq x} \theta(d) \left[\frac{x}{d} \right] = x \sum_{d \leq x} \frac{\theta(d)}{d} + O(x).$$

Let P, Q, T denote the sets of those primes, $\leq x$, for which

$$(2.3) \quad P = \{p; \theta(p) = -1\}, \quad Q = \{q; \theta(q) = 0\}, \quad T = \{t; \theta(t) = 1\}.$$

Let $d(n)$ denote the number of divisors of n .

Then we shall prove

LEMMA 1. With the notations (2.1)–(2.3) we have

$$(2.4) \quad \sum_{n \leq x} g(n) \leq \sum_{n \leq x} d(n) - 2 \sum_{p \in P} \sum_{\substack{p|n \\ n \leq x}} d\left(\frac{n}{p}\right) + 2 \sum_{p \in P} \sum_{\substack{p^2|n \\ n \leq x}} d\left(\frac{n}{p^2}\right) + \\ + 2 \sum_{p' \in P} \sum_{p'' \in P} \sum_{\substack{p'p''|n \\ n \leq x}} d\left(\frac{n}{p'p''}\right) - \sum_{q \in Q} \sum_{q|n} d\left(\frac{n}{q}\right) + 2 \sum_{q \in Q} \sum_{p \in P} \sum_{\substack{pq|n \\ n \leq x}} d\left(\frac{n}{pq}\right).$$

Proof. Let $c(n)d(n)$ be the sum of those terms on the right side which belong to the number n (i.e. the sum of those terms which have the form $d(n/s)$).

We can write n in the form $n = abm$,

$$(2.5) \quad m = p_1^{\alpha_1} \dots p_r^{\alpha_r}, \quad b = q_1^{\beta_1} \dots q_s^{\beta_s}, \quad a = t_1^{\gamma_1} \dots t_u^{\gamma_u}$$

where $p_i \in P, q_i \in Q, t_i \in T$. Then we have

$$(2.6) \quad c(n) = 1 - 2 \sum_{i=1}^r \frac{\alpha_i}{\alpha_i + 1} + 2 \sum_{i=1}^r \frac{\alpha_i - 1}{\alpha_i + 1} + 2 \sum_{i=1}^r \sum_{\substack{j=1 \\ j \neq i}}^r \frac{\alpha_i \alpha_j}{(\alpha_i + 1)(\alpha_j + 1)} - \\ - \sum_{j=1}^s \frac{\beta_j}{\beta_j + 1} + 2 \sum_{j=1}^s \sum_{i=1}^r \frac{\beta_j \alpha_i}{(\beta_j + 1)(\alpha_i + 1)} \\ = 1 + A + B = C + B$$

where

$$(2.7) \quad A = 2 \sum_{i=1}^r \frac{1}{\alpha_i + 1} \left(\alpha_i \sum_{\substack{j=1 \\ j \neq i}}^r \frac{\alpha_j}{\alpha_j + 1} - 1 \right)$$

and

$$(2.8) \quad B = \sum_{j=1}^s \frac{\beta_j}{\beta_j + 1} \left(2 \sum_{i=1}^r \frac{\alpha_i}{\alpha_i + 1} - 1 \right).$$

Now let us regard the following cases:

I. If $r = 0$, i.e. $m = 1$, then

$$c(n) = 1 - \sum_{j=1}^s \frac{\beta_j}{\beta_j + 1} \geq \prod_{j=1}^s \frac{1}{\beta_j + 1} = \frac{1}{d(b)}.$$

If $r \geq 1$ then from (2.7) $B \geq 0$ and so $c(n) \geq C$.

II. If $r = 1$, then

$$C = 1 - \frac{2}{\alpha_1 + 1} = \frac{\alpha_1 - 1}{\alpha_1 + 1} (\geq 0).$$

III. If $r = 2$, and $a_1 = 1$ or $a_2 = 1$, say $a_1 = 1$, then

$$C = 1 - 2 \left(\frac{1}{2} + \frac{1}{a_2 + 1} \right) + 2 \cdot 2 \cdot \frac{1}{2} \cdot \frac{a_2}{a_2 + 1} = \frac{2(a_2 - 1)}{a_2 + 1} \geq 0.$$

IV. If $r = 2$, and $a_i \geq 2$ for $i = 1, 2$ or $r \geq 3$, then for an arbitrary $i \leq r$

$$\alpha_i \sum_{\substack{j=1 \\ j \neq i}}^r \frac{\alpha_j}{\alpha_j + 1} \geq 1$$

and so $A \geq 0$, $C \geq 1$.

On the other hand, one has

$$(2.9) \quad g(n) = \sum_{d|n} \theta(d) = \prod_{p^{\nu} | n} (1 + \theta(p) + \dots + \theta^{\nu}(p)).$$

Hence

$$g(n) = g(a)g(b)g(m) = \begin{cases} d(a) & \text{if } m = l^2, \\ 0 & \text{if } m \neq l^2. \end{cases}$$

So in the following cases we have:

I. $g(n) = d(a) = \frac{d(n)}{d(b)} \leq c(n)d(n);$

II. if a_1 is odd, then $0 = g(n) \leq C \leq c(n) \leq c(n)d(n),$

if a_1 is even, then $g(n) = d(a) \leq \frac{d(n)}{\alpha_1 + 1} \leq Cd(n) \leq c(n)d(n);$

III. $0 = g(n) \leq C \leq c(n) \leq c(n)d(n);$

IV. $g(n) \leq d(a) \leq d(n) \leq Cd(n) \leq c(n)d(n).$

Thus in all cases $g(n) \leq c(n)d(n)$ which proves Lemma 1.

Now let $S = P \cup Q$, and for $s \in S$, let

$$(2.10) \quad \mathbf{s} = \begin{cases} s & \text{if } s \in P, \\ 2s & \text{if } s \in Q. \end{cases}$$

Then using the well-known relations

$$\sum_{\substack{n \leq x \\ l|n}} d\left(\frac{n}{l}\right) = \sum_{m \leq x/l} d(m) = \frac{x}{l} \log \frac{x}{l} + O\left(\frac{x}{l}\right),$$

$$\sum_{\substack{p \leq x \\ p \text{ prime}}} \frac{1}{p} = o(\log x), \quad \sum_{\substack{p \leq x \\ p \text{ prime}}} \sum_{\substack{p' \leq x \\ p' \text{ prime}}} \frac{1}{pp'} = o(\log x),$$

if we add to the right side of (2.4)

$$\frac{1}{2} \sum_{q \in Q} \sum_{\substack{q' \in Q \\ qq' \leq x}} \frac{x}{qq'} \log \frac{x}{qq'} \geq 0,$$

then we have with the notation (2.10) the following

COROLLARY. We have

$$(2.11) \quad \sum_{n \leq x} g(n) \leq x \log x \left(1 - \frac{2}{\log x} U + o(1) \right)$$

where

$$(2.12) \quad U = \sum_{s \in S} \frac{1}{s} \log \frac{x}{s} - \sum_{s \in S} \sum_{\substack{s' \in S \\ ss' \leq x}} \frac{1}{ss'} \log \frac{x}{ss'}.$$

Thus we have to estimate U from below.

Here we shall use the supposition $\sum_{d \leq x} \theta(d) \leq \varepsilon x$ from which

$$\begin{aligned} o(1)x &\geq \sum_{d \leq x} \theta(d) = [x] - 2 \sum_{\substack{d \leq x \\ \theta(d) = -1}} 1 - \sum_{\substack{d \leq x \\ \theta(d) = 0}} 1 \\ &\geq 2 \sum_{p \in P} \frac{x}{p} - \sum_{q \in Q} \frac{x}{q} = x \left(1 - 2 \sum_{s \in S} \frac{1}{s} \right) - 1 \end{aligned}$$

follows. Hence

$$(2.13) \quad \sum_{s \in S} \frac{1}{s} \geq \frac{1}{2} - o(1).$$

Let

$$S = \{s_1 < s_2 < \dots < s_k\}$$

and

$$(2.14) \quad S' = \left\{ s_1 < s_2 < \dots < s_l; \sum_{i=1}^{l+1} \frac{1}{s_i} > \frac{1}{2} \geq \sum_{i=1}^l \frac{1}{s_i} \right\}.$$

(If $\sum_{s \in S} \frac{1}{s} \leq \frac{1}{2}$ then let $S' = S$.)

Now we define a $\theta'(n)$ completely multiplicative function for $n \leq x$ with

$$\begin{aligned} \theta'(p) &= -1 & \text{if } p \in P \cap S', \\ \theta'(q) &= 0 & \text{if } q \in Q \cap S', \\ \theta'(t) &= 1 & \text{if } t \in T \cup (S \setminus S'). \end{aligned}$$

Thus we have from (2.9) for an arbitrary n

$$(2.15) \quad \sum_{d|n} \theta(d) = g(n) \leq g'(n) = \sum_{d|n} \theta'(d).$$

So we shall use the Corollary for $\theta'(n)$ and we shall estimate the corresponding U' , i.e. we shall prove

LEMMA 2. We have

$$(2.16) \quad U' = \sum_s \frac{1}{s} \log \frac{x}{s} - \sum_s \sum_{s' \leq x/s} \frac{1}{ss'} \log \frac{x}{ss'} \geq \left(\frac{1}{\sqrt{e}} - \frac{1}{2} + o(1) \right) \log x$$

where the summation in Lemma 2 runs always through $s, s' \in S'$.

Proof. If in the definition of S' in (2.14) $s_{i+1} \leq \log x$, and $s \neq s$ then as $\sum_s \frac{1}{s} > \frac{1}{6}$, we get

$$\begin{aligned} U' &\geq \sum_s \frac{1}{s} \log x (1 - o(1)) - \log x \sum_s \sum_{s'} \frac{1}{ss'} \\ &\geq (1 - o(1)) \log x \left(\sum_s \frac{1}{s} \right) \left(1 - \sum_{s'} \frac{1}{s'} \right) \geq \left(\frac{1}{6} \cdot \frac{5}{6} - o(1) \right) \log x \end{aligned}$$

which implies (2.16).

If $s_{i+1} > \log x$, or $s = s'$ let

$$(2.17) \quad \alpha = \sum_{s \leq \sqrt{x}} \frac{1}{s}, \quad \beta = \sum_{\sqrt{x} < s \leq x} \frac{1}{s}.$$

Then as $1/s_{i+1} = o(1)$, we have

$$(2.18) \quad \frac{1}{2} - o(1) \leq \alpha + \beta \leq \frac{1}{2}.$$

With these notations the following formulae hold

$$\begin{aligned} (2.19) \quad D &= \sum_{s \leq \sqrt{x}} \frac{1}{s} \log \frac{x}{s} - \sum_{s \leq \sqrt{x}} \sum_{s' \leq \sqrt{x}} \frac{1}{ss'} \log \frac{x}{ss'} \\ &\geq \sum_{s \leq \sqrt{x}} \frac{1}{s} \log \frac{x}{s} \left(1 - \sum_{s' \leq \sqrt{x}} \frac{1}{s'} \right) \\ &= (1 - \alpha) \log x \cdot \alpha - (1 - \alpha) \sum_{s \leq \sqrt{x}} \frac{\log s}{s}, \end{aligned}$$

$$\begin{aligned} (2.20) \quad E &= \sum_{s \leq \sqrt{x}} \frac{1}{s} \sum_{\sqrt{x} < s' \leq x/s} \frac{1}{s'} \log \frac{x}{ss'} \leq \sum_{\sqrt{x} < s' \leq x} \frac{1}{s'} \sum_{s \leq \sqrt{x}} \frac{1}{s} \log \frac{\sqrt{x}}{s} \\ &= \frac{1}{2} \beta \log x \cdot \alpha - \beta \sum_{s \leq \sqrt{x}} \frac{\log s}{s}, \end{aligned}$$

$$(2.21) \quad F = \sum_{s > \sqrt{x}} \frac{1}{s} \log \frac{x}{s} = \beta \log x - \sum_{\sqrt{x} < s \leq x} \frac{\log s}{s}.$$

Here as $U' = D - 2E + F$, from formulae (2.18)–(2.21) we get

$$\begin{aligned} (2.22) \quad U' &\geq \log x [\alpha(1 - \alpha - \beta) + \beta] - (1 - \alpha - 2\beta) \sum_{s \leq \sqrt{x}} \frac{\log s}{s} - \sum_{\sqrt{x} < s \leq x} \frac{\log s}{s} \\ &= \left(\frac{1}{2} - \frac{\alpha}{2} + o(1) \right) \log x - (\alpha + o(1)) \sum_{s \leq \sqrt{x}} \frac{\log s}{s} - \sum_{\sqrt{x} < s \leq x} \frac{\log s}{s}. \end{aligned}$$

On the other hand, it is easy to show that if S' is a set of primes $s \leq y$ ($y \neq O(1)$), $s = s$ or $2s$, and

$$(2.23) \quad \sum_{s \in S'} \frac{1}{s} = \gamma + o(1) \quad (\leq 1)$$

(γ a given number) then the sum

$$\sum_{s \in S'} \frac{1}{s}$$

is maximal if the set S' contains all primes in an interval $[z, y]$ and no primes less than z and for all primes s , $s = s$. So if we use the formulae

$$\sum_{\substack{p \leq x \\ p \text{ prime}}} \frac{\log p}{p} = \log x + O(1), \quad \sum_{\substack{p \leq x \\ p \text{ prime}}} \frac{1}{p} = \log \log x + C + o(1)$$

where C is an absolute constant, easy computation shows that if (2.23) holds then

$$(2.24) \quad \sum_{s \in S'} \frac{\log s}{s} \leq (1 - e^{-\gamma} + o(1)) \log y.$$

Thus from (2.22) and (2.24) we have

$$(2.25) \quad U' \geq \log x \left(\frac{1}{2} - \frac{\alpha}{2} - \frac{\alpha}{2} (1 - e^{-\alpha}) - (1 - e^{-\beta}) - o(1) \right) = G \cdot \log x.$$

Here using $\beta = \frac{1}{2} - a + o(1)$, we get with some computation that for $0 \leq a \leq \frac{1}{2} + o(1)$

$$(2.26) \quad G = G(a, \beta) = G(a) \geq G(0) = \frac{1}{\sqrt{e}} - \frac{1}{2} - o(1),$$

which proves Lemma 2.

Thus we have from formulae (2.2), (2.15), (2.11), (2.12) and (2.16)

$$(2.27) \quad \sum_{d \leq x} \frac{\theta(d)}{d} = \frac{1}{x} \sum_{n \leq x} g(n) + O(1) \leq \frac{1}{x} \sum_{n \leq x} g'(n) + O(1) \\ \leq \log x \left(1 - \frac{2U'}{\log x} + o(1) \right) \leq \log x \left(2 \left(1 - \frac{1}{\sqrt{e}} \right) + o(1) \right).$$

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**The factorization of $Q(L(x_1), \dots, L(x_k))$
over a finite field where $Q(x_1, \dots, x_k)$
is of first degree and $L(x)$ is linear**

by

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I. Introduction. Let $\text{GF}(q)$ denote the finite field of order $q = p^n$ where p is prime and $n \geq 1$. Let $\Gamma(p)$ denote the algebraic closure of $\text{GF}(p)$. A polynomial $Q \in \text{GF}[q; x_1, \dots, x_k]$ is *absolutely irreducible* if Q has no nontrivial factors over $\Gamma(p)$. Throughout this paper, the term irreducible will mean absolutely irreducible.

A polynomial with coefficients in $\text{GF}(q)$ of the form

$$L(x) = \sum_{i=0}^r c_i x^i$$

is called a linear polynomial. The requirement that the coefficients be in $\text{GF}(q)$ insures that the operation of mapping composition for linear polynomials is commutative. Corresponding to the linear polynomial $L(x)$ we have the ordinary polynomial

$$l(x) = \sum_{i=0}^r c_i x^i.$$

We shall assume in the following that $c_0 \neq 0$; this avoids multiple factors in $L(x)$ and insures that there is a smallest integer r such that $l(x)$ divides $x^r - 1$. We say that $l(x)$ has *exponent* r .

Let $Q(x_1, \dots, x_k) = a_1 x_1 + \dots + a_k x_k + 1$ where $[\text{deg } a_1, \dots, \text{deg } a_k] = s$ (if $a \in \text{GF}(q^s)$ but $a \notin \text{GF}(q^t)$, $1 \leq t < s$, we say that the *degree of a relative to $\text{GF}(q)$* is s and write $\text{deg } a = s$). We shall assume that $\{a_1, \dots, a_k\}$ are linearly independent over $\text{GF}(q)$; otherwise $Q(x_1, \dots, x_k)$ can be reduced at once to a polynomial in m variables by suitable first degree transformations, where m is the number of elements in a maximal linearly independent subset of a_1, \dots, a_k .

In this paper we describe the factorization of $Q(L(x_1), \dots, L(x_k))$. (We note that it is possible to have $Q(L(x_1), \dots, L(x_k))$ reduce to a polynomial in fewer than k variables even though $\{a_1, \dots, a_k\}$ are linearly